

SH6610 instructions introduction

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Analysis of SH6610 Instructions

The following a re SH6610 in structions, catego rized a nd ex plained acco rding to their respective functions. When you need an instruct ion for a certain function, you can look up the instruction in its function category. You'd better browse through all the instructions listed here, because even thoug h you can't rememb er all of them at once, you can still have an impression to remind you of such inst ruction when need ed. Some people don 't even know about some very good instructions that he can use, just because he didn't browse through all of them. This ignorance may cause considerable waste of space on a system without a b ig enough ROM, which is so regrettable. I will ex plain as simply as possible to improve your learning efficiency. Of co urse, it is n ormal to u nderstand or remember only part of the contents after reading on ce. And you are certain often revert to this bo ok because it is a collection of instructions.

Instructions

Instruction is a se ries of <u>codes</u> that can be recognized by CPU, and then CPU will operate according to the given instruction.

Operand

Apart from instructions t o tell CPU how to op erate, the ope ration o bject must be designated. The o bject for CP U to o perate with is called <u>Operand</u>. Therefore, a complete instruction must include two items, both <u>instruction</u> and <u>operand</u>.

Format of SH6610 Instructions

Instruction [Operand 1], [Operand 2]

Of the above, items inside [] are used according to nature of the instruction. Some instructions require only one operand, while some require two. Instruction and operand shall be separated by a space, and operands shall be separated from each other by a ','.

Execution Time of Instruction

Execution time for SH6610 instructions is <u>one instruction cycle</u>, which is one fourth of the system working frequency.



1-1 Instruction Categories

According to their functions, SH6610 instructions can be classified into four categories:

- Arithmetic Operation Instructions
 - I. addition: ADC , ADCM , ADD , ADDM , ADI , ADIM
 - II. Subtraction: SBC , SBCM , SUB , SUBM , SBI , SBIM
 - III. BCD: DAA , DAS
- Logic Operation Instructions
 EOR , EORM , EORIM , OR , ORM , ORIM , AND , ANDM , ANDIM
- Data Transmission Instructions

LDA, STA, LDI

• Flow Control Instructions

BAZ , BC , BA0 , BA1 , BA2 , BA3 , CALL , RTNW , RTNI , HALT , STOP , JMP , TJMP

The above a re all of the instructions of SH6610 series, which a dd up to only 40 in number, but never overlook them! A variety of consumer electrical products on the market are created with them, e.g. calculator, remote controller, watch, toy, etc.

1-2 Explanation of Symbols

Before going to our subject, we list the symbols that may appear afterwards so that our readers can understand this book more easily. The symbols are listed as follows:

PC	Program Counter	
AC	Accumulator	
CY	Carry Flag	
Mx	Data Memory	
bbb RAM	bank	
ST	Stack	
TBR	Table Branch Register	
Х	Program Address	
I	Immediate Data	



& Logic AND

- Logic OR
- ^ Logic EOR

Now let's enter the world of instructions of SH6610 series. "Let 's go $\ensuremath{^\circ_\circ}$



1-3 Instructions for Data Transmission

During the internal o peration of the system, da ta is transmit ted rapidly a nd incessantly between memories or registers. This fast and incessant transmission is the power of system capability. SH6610 system provides several instructions for data transmission, as follows:

Instruction: LDI	Function: to load Immediate Data I to "Accumulator" and "Data Memory"	
Format:	LDI Mx , I	
Instruction Code:	01111 iiii xxx xxxx	
Carry Flag:	Not affected	
Operation:	AC , Mx ← I	

Explanation

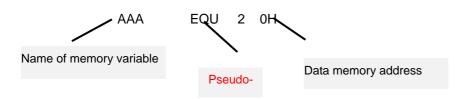
LDI is a very frequently used inst ruction. It loads immediate d ata I to accu mulator and data memory. However, due to this 4-bit system, the preset range of the immediate data I is $00H \sim 0FH(0 \sim 15)$, and that of Mx is $00H \sim 7FH$.

[Example]

LDI 20H , 05H after execution: AC=05H Content of data memory \$20H=05H

Programming Tip

There isn't any specially defined register for users in SH6610 system, but we can use its powerful data memory as registers in our program designing. In the above example, I have used data memory \$20H for a register. However, if they are expressed only by address and without respective names, the design of the program will be very confusing. Here's a tip for you: you can use the pseudo-instruction EQU to define each data address.



Therefore, the LDI instruction can also be written like this: LDI AAA , 05H



Instruction: STA	Function:to store the value of Accumulator to Data Memory
Format:	STA Mx , bbb
Instruction Code: Carry Flag: Operation:	00111 1bbb xxx xxxx Not affected Mx ← AC

STA loads the value of accumulator to dat a memory. When executing this instruction, CPU does the transmission only and the carry flag is not affected.

[Example] Save the value of AC in \$21H

	:	
LDI	20H , 05H	;AC=05H , \$20H=05H
STA	21H , 00H	;\$21H=05H
	:	

Programming Tip

If operand 2 is the immediate data I, then the program can be written in the following ways:

LDI	20H, 0AH	;expressed in hex
LDI	20H, 10	;expressed in decimal
LDI	20H, 1010B	,expressed in binary

Instruction: LDA	Function: to load the value of Data Memory to Accumulator
Format:	LDA Mx , bbb
Instruction Code:	00111 0bbb xxx xxxx
Carry Flag:	Not affected
Operation:	$AC \leftarrow Mx$

Explanation

LDA loads the value of da ta memory to accumulator. This instruction does not affect the carry flag.



[Example] Load the value of \$20H to AC

		:	
	LDI	20H,05H	;\$20H=05H,AC=05H
LDI		21H,0FH	;\$21H=0FH,AC=0FH
LDA		20H,0	;AC=05H
		:	

Programming Tip

When writing a prog ram, we often use la bels as jumping de stinations in the program. Label names can be defined by user according to the following rules:

- I. A label mustn't begin with number or space.
- *II.* Length of a label mustn't exceed 7 characters. Only the first 7 characters will be recognized for labels exceeding that length.



1-4 Instructions for Arithmetic Operation

SH6610 series provide some frequently used arithmetic (integer) operation instructions, like: Addition, subtraction, BCD adjustment, etc.

Instruction: ADD	Function: to add up the values of Data Memory and Accumulator,	
	and then save the result in the Accumulator	
Format:	ADD Mx , bbb	
Instruction Code:	00001 0bbb xxx xxxx	
Carry Flag:	CY	
Operation:	$AC \leftarrow Mx + AC$	

Explanation

Instruction ADD adds up the values of data memory and accumulator and saves the result in the accumulator. The ADD operation affects carry flag: when the result of ADD exceeds 0FH, the carry flag is set to 1; otherwise the value of CY is 0. Therefore we can decide whether there's a carry by the value of the carry flag after addition.

[Example] 05H +06H

1

LDI	20H,05H	;\$20H=05H,AC=05H
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=05H
ADD	21H,0	;AC=0BH,CY=0,\$21H=06H
		:

[Example] 0BH + 06H

	•	
LDI	20H,0BH	;\$20H=0BH,AC=0BH
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=0BH,CY=0
ADD	21H,0	;AC=01H,CY=1,\$21H=06H

Instruction: ADDM	Function: to add up the values of Data Memory and Accumulator, and save the result in both the Accumulator and the Data Memory
Format:	ADDM Mx , bbb
Instruction Code:	00001 1bbb xxx xxxx
Carry Flag:	CY
Operation:	AC , $Mx \leftarrow Mx + AC$



Instruction ADDM a dds up the valu es of data memory and accumulator and saves the result in both the accumulator and the data memory. The ADDM operation affects carry flag: when the result of ADDM exceeds 0FH, the carry flag is set to 1; otherwise the value CY is 0. Therefore we can decide whether there's a carry by the value of the carry flag after addition.

[Example] 05H +06H

:

LDI	20H,05H	;\$20H=05H,AC=05H
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=05H
ADDM	21H,0	;AC=0BH,\$21H=0BH,CY=0
		:

[Example] 0BH + 06H

	•	
LDI	20H,0BH	;\$20H=0BH,AC=0BH
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=0BH,CY=0
ADDM	21H,0	;AC=01H,\$21H=01H,CY=1

Programming Tip

When you are reading the examples, I suggest that you call them into ICE after compiling to watch the change in each of the registers step by step.

Instruction: ADC	Function: to add up the value of Data Memory, Carry Flay and the value of Accumulator, and then save the result in the Accumulator
Format:	ADC Mx , bbb
Instruction Code:	00000 0bbb xxx xxxx
Carry Flag:	CY
Operation:	$AC \leftarrow Mx + AC + CY$

Explanation

Instruction ADC ad ds up the value of dat a memory, carry flag and the value of



accumulator, and saves the result in the accumulator. The ADC operation affects the carry flag: when the result of ADC exceeds 0FH, the carry flag is set to 1; otherwise the value of CY is 0. Therefore we can d ecide whether there's a carry by the value of the carry flag after addition.

```
[Example] 05H +06H, CY=1
```

	:	;CY=1
LDI	20H,05H	;\$20H=05H,AC=05H
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=05H
ADC	21H,0	;AC=0CH,\$21H=06H,CY=0
	:	

```
[Example] 0BH + 06H , CY=0
```

	:	;CY=0
LDI	20H,0BH	;\$20H=0BH,AC=0BH
LDI	21H,06H	;\$21H=06H,AC=06H
LDA	20H,00H	;AC=0BH
ADC	21H,0	;AC=01H,\$21H=06H,CY=1
	:	

Programming Tip

When carry is not considered, you'd better use ADD rather than ADC, in order to a void extra uncertainty due to the addition of CY value (because CY can be either 1 or 0).



Instruction: ADCM	Function: to add up the value of Data Memory, Carry Flag and the value of Accumulator, and then save the result in both the	
	Accumulator and the Data Memory	
Format:	ADCM Mx , bbb	
Instruction Code:	00000 1bbb xxx xxxx	
Carry Flag:	CY	
Operation:	AC , $Mx \leftarrow Mx + AC + CY$	

Instruction ADCM ad ds up the value of dat a memory, ca rry flag and the value of accumulator, and saves the result in both the accumulator and the data memory. The ADCM operation affects the carry flag: when the result of ADCM exceeds 0FH, the carry flag is set to1; otherwise the value of CY is 0. Therefore we can decide whether there's a carry by the value of the carry flag after addition.

[Example] 05H +06H , CY=1

	: ;CY=1			
LDI	20H,05H	;\$20H=05H,AC=05H		
LDI	21H,06H	;\$21H=06H,AC=06H		
LDA	20H,00H	;AC=05H		
ADCM	21H,0	;AC=0CH,\$21H=0CH,CY=0		
	:			

```
[Example] 0BH + 06H , CY=0
```

	: ;CY=0			
LDI	20H,0BH	;\$20H=0BH,AC=0BH		
LDI	21H,06H	;\$21H=06H,AC=06H		
LDA	20H,00H	;AC=0BH,		
ADCM	21H,0	;AC=01H,\$21H=01H,CY=1		

Instruction: ADI	Function: to add up the value of Data Memory and Immediate Data I, and then save the result in Accumulator
Format:	ADI Mx , I
Instruction Code:	01000 iiii xxx xxxx
Carry Flag:	CY
Operation:	AC ← Mx + I

Explanation

Instruction ADI adds up the value of dat a memory and immediate data I, and save s the



result in accumulator. The ADI operation affects carry flag: when the result of ADI exceed s 0FH, the carry flag is set to1; otherwi se the value of CY is 0. Therefore we can decide whether there's a carry by the value of the carry flag after addition.

[Example] \$20H=0 5H, I=04H

:

LDI	20H,05H	;\$20H=05H,AC=05H,CY=0
ADI	20H,04H	;AC=09H,\$20H=05H,CY=0
	:	

[Example] \$20H=0 AH , I=07H

	:	
LDI	20H,0AH	;\$20H=0AH,AC=0AH,CY=0
ADI	20H,07H	;AC=01H,\$20H=0AH,CY=1

Instruction: ADIM	Function: to add up the value of Data Memory and Immediate Data I, and then save the result in both Accumulator and the Data Memory
Format:	ADIM Mx , I
Instruction Code:	01001 iiii xxx xxxx
Carry Flag:	CY
Operation:	AC , $Mx \leftarrow Mx + I$

Explanation

Instruction ADIM adds up the value of data memory and immediate dat a I, and saves the result in both accumulator and the dat a memory. The ADIM o peration affects carry flag: when the result of ADIM exceeds 0FH, the carry flag is set to1; otherwise the value of CY is 0. Therefore we can decide by the value of the carry flag after addition whether there is a carry.

[Example] \$20H=0 5H, I=04H

	:		
	LDI	20H,05H	;\$20H=05H,AC=05H
	ADIM	20H,04H	;AC=09H,\$20H=09H,CY=0
	:		
[Example] \$20H=0) AH , I=07H		
	:		
	LDI	20H,0AH	;\$20H=0AH,AC=0AH



ADIM

:

20H,07H ;AC=01H,\$20H=01H,CY=1

Instruction: DAA	Function: to adjust the value of Data Memory to decimal after	
	addition, and then save the result in both Accumulator	
	and the Data Memory	
Format:	DAA Mx	
Instruction Code:	11001 0110 xxx xxxx	
Carry Flag:	CY	
Operation:	AC ; Mx ← Decimal , adjust AC for Add	

Explanation

Instruction DAA acts by adjusting the value of data memory to decimal after addition and saving the re sult to both accumul ator and the <u>data memory</u>. Its adjusting m ethod is if the value of the data memory is greater than 9 or if CY= 1, then add 6 to the data memory and set the carry flag to 1.

[Example] 06H + 05H , and do DAA adjustment

	:	
LDI	20H,06H	;AC=06H,\$20H=06H
LDI	21H,05H	;AC=05H,\$21H=05H
LDA	20H,0	;AC=06H
ADD	21H,0	;AC=0BH,CY=0
DAA	21H	;AC=01H,\$21H=01H,CY=1
	:	



Instruction: SUB	Function: to subtract the value of Accumulator from the value of Data Memory, and then save the result in the Accumulator
Format:	SUB Mx , bbb
Instruction Code:	00011 0bbb xxx xxxx
Carry Flag:	CY
Operation:	AC ← Mx - AC

SUB subtracts the value of accumulator from the value of data memory and saves the result in the accumulator. When executing SUB, if the value of data memory is less than the value of accumulator, a "borrow" will take place and CY will be set to 0. On the contrary, if the value of dat a memory is gre ater than the value of accumulator, borrow will not happe n and CY will be set to 1. Therefore we can decide by the value of CY whether there is a borrow after execution of SUB. Besides, subtraction in the system is done through addition, i.e. when subtracting a number, it is actually adding the number's binary complement.

[Example] 06H - 05H

:

:

LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	20H,0	;AC=05H
SUB	21H,0	;AC=1,CY=1,\$21H=06H
	:	

[Example] 05H - 06H

LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	21H,0	;AC=06H
SUB	20H,0	;AC=0FH,CY=0,\$20H=05H
	:	



Instruction: SUBM	Function: to subtract the value of Accumulator from the value of Data Memory, and then save the result in both the
	Accumulator and the Data Memory
Format:	SUBM Mx , bbb
Instruction Code:	00011 1bbb xxx xxxx
Carry Flag:	CY
Operation:	AC , $Mx \leftarrow Mx - AC$

System movement of SUBM is alm ost the sam e as SUB, it subtracts current value of accumulator from the value of data memory and saves the result in the accumulator as well as in the <u>data memory</u>. When executing SUBM, if the value of data memory is less than the value of accumulator, a "borrow" will take place and CY will be set to 0. On the contrary, if the value of dat a memory is gre ater than the value of accumulator, borrow will not happe n and CY will be 1. Therefore we can decide by the value of CY whether there is a borr ow after execution of SUBM.

[Example] 06H - 05H

:

:

LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	20H,0	;AC=05H
SUBM	21H,0	;AC=01H,CY=1,\$21H=01H
	:	

[Example] 05H - 06H

LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	21H,0	;AC=06H
SUBM	20H,0	;AC=0FH,CY=0,\$20H=0FH
	:	



Instruction: SBC	Function: to subtract the value of Accumulator from the value of Data Memory, add Carry Flag, and then save the result in
	the Accumulator
Format:	SBC Mx , bbb
Instruction Code:	00010 0bbb xxx xxxx
Carry Flag:	CY
Operation:	$AC \leftarrow Mx - AC + CY$

System movement of SB C is to subtract the value of accum ulator from the value of data memory, add the value of carry flag, and then save the result in the accumulator. When executing SBC, if the value of data memory is less than the value of accumulator, a "borrow" will take place and the CY will be set to 0. On the contrary, if the value of dat a memory is greater than the value of accum ulator, bo rrow will not hap pen and the CY will be 1. Therefore we can decide by the value of CY whether there is a borrow after execution of SBC.

[Example] CY=0 , 6 - 5=?

	:	;CY=0
LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	20H,0	;AC=05H
SBC	21H,0	;AC=01H,CY=0,\$21H=06H
	:	

[Example] CY=1 , 6 - 5=?

	:	;CY=1
LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	20H,0	;AC=05H
SBC	21H,0	;AC=02H,CY=0,\$21H=06H
	:	

Instruction: SBCM	Function: to subtract the value of Accumulator from the value of Data Memory, add Carry Flag, and then save the result in	
	both the Accumulator and the Data Memory	
Format:	SBCM Mx , bbb	
Instruction Code:	00010 1bbb xxx xxxx	
Carry Flag:	CY	
Operation:	AC , $Mx \leftarrow Mx - AC + CY$	



System movement of SBCM is to subtract the value of accumulator from the value of data memory, add the value of carry flag, and then save the result in both the accumulator and <u>data memory</u>. When executing SBCM, if the value of d ata memory is I ess than the value of accumulator, a "borrow" will take place and the CY will be set to 0. On the contrary, if the value of data memory is greater t han the value of accumulator, borrow will not hap pen and the CY will be set to 1. Therefore we can decide by the value of CY whether there is a borrow after execution of SBCM.

```
[Example]CY=0, 6-5=?
```

	:	;CY=0
LDI	20H,05H	;AC=05H,\$20H=05H
LDI	21H,06H	;AC=06H,\$21H=06H
LDA	20H,0	;AC=05H,CY=0
SBCM	21H,0	;AC=01H,CY=0,\$21H=01H
	:	

[Example]CY=1, 6-5=?

	:	;CY=1	
LDI	20H,05H	;AC=05H,	\$20H=05H
LDI	21H,06H	;AC=06H,	\$21H=06H
LDA		20H,0	;AC=05H,CY=1
SBCM		21H,0	;AC=02H,CY=0,\$21H=02H
	:		

Instruction: SBI	Function: to subtract the Immediate Data I from the value of Data Memory, and then save the result in Accumulator
Format:	SBI Mx. I
Instruction Code:	01010 iiii xxx xxxx
Carry Flag:	CY
Operation:	AC ←Mx - I

Explanation

System movement of SBI is to subtract immediate data I from the value of data memory, and save the result in accumulator. When executing SBI, if the value of the dat a memory is less than the immediate d ata, a "bo rrow" will t ake place and CY will be set to 0. On the



contrary, if the value of dat a memory is gr eater than the immediate data, borrow will not happen and CY will be 1. Therefore we can decide by the value of CY whether there is a borrow after execution of SBI.

[Example] \$20H=0 5H, I=04H

:

		•	
	LDI	20H,05H	;\$20H=05H,AC=05H,CY=0
	SBI	20H,04H	;AC=01H,\$20H=05H,CY=1
		:	
[Examp	le] \$20H	=0 2H , I=07H	
[Examp	le] \$20H:	=0 2H , I=07H :	
[Examp	le] \$20H: LDI	=0 2H , I=07H : 20H,02H	;\$20H=02H,AC=02H,CY=0
[Examp	-	:	;\$20H=02H,AC=02H,CY=0 ;AC=0BH,\$20H=02H,CY=0

Instruction: SBIM	Function: to subtract Immediate Data I from the value of Data Memory, and then save the result in both Accumulator	
	and the Data Memory	
Format:	SBIM Mx , I	
Instruction Code:	01011 iiii xxx xxxx	
Carry Flag:	CY	
Operation:	AC , $Mx \leftarrow Mx - I$	

Explanation

System movement of SBIM is to subtract immediate data I from the value of data memory, and save the result in both accumulator and the <u>data memory</u>. When executing SBIM, if the value of data memory is less than the immediate data, a "borrow" will take place and CY will be set to 0. On the cont rary, if the value of dat a memory is g reater than the immediate data, borrow will not happen and CY will be 1. Therefore we can decide by the value of CY whether there is a borrow after execution of SBIM.

[Example] \$20H=0 5H, I=04H

LDI	20H,05H	;\$20H=05H,AC=05H
SBIM	20H,04H	;AC=01H,\$20H=01H,CY=1
	:	



[Example] \$20H=0 2H, I=07H

:

:

LDI	20H,02H	;\$20H=02H,AC=02H
SBIM	20H,07H	;AC=0BH,\$20H=0BH,CY=0

Instruction: DAS	Function: to adjust the value of Data Memory to decimal after	
	subtraction, and save the result in Accumulator and the	
	Data Memory	
Format:	DAS Mx	
Instruction Code:	11001 1010 xxx xxxx	
Carry Flag:	CY	
Operation:	AC ; Mx ← Decimal , adjust AC for Sub	

Explanation

Instruction DAS acts by adjusting the value of data memory to decimal after subtraction and saving the result to b oth accumulator and the data memory. Its adjusting method is if the value of data memory is greater than 9 or if CY=0, then add 0AH to the data memory and set CY to 0.

[Example] 05H - 06H , and do DAS adjustment

	:	
LDI	20H,06H	;AC=06H,\$20H=06H
LDI	21H,05H	;AC=05H,\$21H=05H
LDA	20H,0	; AC=06H
SUB	21H,0	;AC=0FH,CY=0
AS	21H	;AC=09H,\$21H=09H,CY=0
	:	

D



1-5 Instructions for Logic Operation

Logic instructions are essential to system structure. SH6610 se ries MCU provid e some common logic instructions. Now I'm going to explain to you one by one in most details, and assist my explanation with simple examples, so that you can quickly understand action theory of each instruction.

Instruction: AND	Function: to do logic AND operation with the values of Data Memory	
	and Accumulator, and then save the result in the	
	Accumulator	
Format:	AND Mx , bbb	
Instruction Code:	00110 1bbb xxx xxxx	
Carry Flag:	Not affected	
Operation:	$AC \leftarrow Mx \& AC$	

Explanation

In AND operation, the result will be 1(true) only if both of the two operands are 1(true). Its logic table is as follows:

Logic operation table for AND

а	b	a AND b
0	0	0
0	1	0
1	0	0
1	1	1

However, in real inst ruction the logic operand has 4 bit s rather than 1 bit. The instruction A ND is to AND the values of dat a memory with accumulator, and the re sult is saved in the accumulator.

[Example] 06H & 05H

2

LDI	20H,0110B	;AC=06H,\$20H=06H
LDI	21H,0101B	;AC=05H,\$21H=05H
AND	20H,0	;AC=0100B,\$20H=06H
:		



Instruction: ANDM	Function: to do logic AND operation with the values of Data
	Memory and Accumulator, and then save the result in
	both the Accumulator and the Data Memory
Format:	ANDM Mx , bbb
Instruction Code:	00110 1bbb xxx xxxx
Carry Flag:	Not affected
Operation:	AC , Mx ← Mx & AC

Explanation

System movement of the instruction ANDM is almost the same as AND, but saving the operation result in <u>data memory</u> as well as in accumulator.

[Example] 01 10B & 0101B

:

	:	
LDI	20H,0110B	;AC=0110B,\$20H=0110B
LDI	21H,0101B	;AC=0101B,\$21H=0101B
ANDM	20H,0	;AC=0100B,\$20H=0100B

Instruction: ANDIM	Function: to do logic AND operation with the value of Data Memory and Immediate Data I, and then save the result in both Accumulator and the Data Memory
Format:	ANDIM Mx , I
Instruction Code:	01110 iiii xxx xxxx
Carry Flag:	Not affected
Operation:	AC, $Mx \leftarrow Mx \& I$

Explanation

System movement of the instruction ANDIM is to change operand 2 (accumulator) of instruction AND to imme diate data I. This instruction is in imme diate mode, so the add ress of data memory can only be set to bank 0($000 \text{ H} \sim 07\text{FH}$). The operation result is saved in both accumulator and the data memory.

[Example] \$20H=01 10B, I=0011B

:

LDI	20H,0110B	;AC=06H,\$20H=0110B
ANDIM	20H,0011B	;AC=0010B,\$20H=0010B
	:	



Programming Tip

ANDIM itself has a special function "MASK". When we need to set a certain bit to 0, we can clear this bit to 0 with ANDIM like this:

[Example] Clear bit 2 of \$20H to 0

ANDIM 20H , 1011B

Af ter execution: \$20H=x0xxB



Instruction: OR	Function: to do logic OR operation with the values of Data Memory and Accumulator, and then save the result in the
	Accumulator
Format:	OR Mx , bbb
Instruction Code:	00101 0bbb xxx xxxx
Carry Flag:	Not affected
Operation:	$AC \leftarrow Mx AC$

In OR operation, the result will be 1(true) if either one of the two operands is 1(true). Its logic table is as follows:

Logic operation table for OR

а	b	a OR b
0	0	0
0	1	1
1	0	1
1	1	1

However, in real inst ruction the l ogic op erand also h as 4bits rath er than 1 bit. The instruction OR is to OR the values of data memory with accumulator, and the result is saved in the accumulator.

[Example] 0001B | 0100B

	:	
LDI	20H,0001B	;\$20H=0001B,AC=0001B
LDI	21H,0100B	;\$21H=0100B,AC=0100B
OR	20H,0	;\$20H=0001B,AC=0101B
	:	

Instruction: ORM	Function: to do logic OR operation with the values of Data Memory and Accumulator, and then save the result in both the	
	Accumulator and the Data Memory	
Format:	ORM Mx , bbb	
Instruction Code:	00101 1bbb xxx xxxx	
Carry Flag:	Not affected	
Operation:	AC , Mx← Mx AC	



System movement of the in struction O RM is almost the same a s O R, but saving t he operation result in <u>data memory</u> as well as in accumulator.

[Example] 0001B | 0100B

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LDI	20H,0001B	;\$20H=0001B,AC=0001B
LDI	21H,0100B	;\$21H=0100B,AC=0100B
ORM	20H,0	;\$20H=0101B,AC=0101B
	:	

Programming Tip

In program designing, if a cert ain bit of a variable needs to be set to 1 and the other bits must not be affected, then this can be done by the instruction OR. Because any bit that has done OR operation with 0 can keep its original value, while those with 1 will have the value of 1.

Instruction: ORIM	Function: to do logic OR operation with the value of Data Memory		
	and Immediate Data I, and then save the result in both		
	Accumulator and the Data Memory		
Format:	ORIM Mx , I		
Instruction Code:	01101 iiii xxx xxxx		
Carry Flag:	Not affected		
Operation:	AC, $Mx \leftarrow Mx \mid I$		

Explanation

System movement of the instruction ORIM is to do logic OR operation with the value of data memory and Imm ediate Data I, and save the result in b oth accumulator and the data memory. This instruction is also in immediate mode.

[Example] Set bit 3 of the value of \$20H to 1

ORIM 20H , 1000B After execution: \$20H=1xxxB



Instruction: EOR	Function: to do logic Exclusive OR operation with the values of Data Memory and Accumulator, and then save the result in the Accumulator
Format:	EOR Mx , bbb
Instruction Code:	00100 0bbb xxx xxxx
Carry Flag:	Not affected
Operation:	AC ← Mx ^ AC

System movement of EOR is to do logic Exclu sive OR operation with the value s of data memory and accumulator, and save the result in the accumulator. EOR is usually referred to as Exclusive OR, becau se the operati on result will be 1 only if the two opera nds have different values; otherwise the result will be 0. The logic table for EOR is as follows:

Logic operation table for EOR

а	b	a EOR b
0	0	0
0	1	1
1	0	1
1	1	0

[Example] 0011B ^ 0101B

.

	•	
LDI	20H,0011B	;\$20H=0011B,AC=0011B
LDI	21H,0101B	;\$21H=0101B,AC=0101B
EOR	20H,0	;\$20H=0011B,AC=0110B
	:	



Instruction: EORM	Function: to do logic Exclusive OR operation with the values of		
	Data Memory and Accumulator, and then save the result		
	in both the Accumulator and the Data Memory		
Format:	EORM Mx , bbb		
Instruction Code:	00100 1bbb xxx xxxx		
Carry Flag:	Not affected		
Operation:	AC , $Mx \leftarrow Mx \wedge AC$		

System movement of the instruction EORM is almost the same as EOR, i.e. doing EOR action with the values of data memory and accumulator, but saving the operation result in the <u>data memory</u> as well as in the accumulator.

[Example] 0011B ^ 0101B

	•	
LDI	20H,0011B	;\$20H=0011B,AC=0011B
LDI	21H,0101B	;\$21H=0101B,AC=0101B
EORM	20H,0	;\$20H=0110B,AC=0110B
	:	

Instruction: EORIM	Function: to do logic EOR operation with the value of Data Memory and Immediate Data I, and then save the result in both	
	the Accumulator and the Data Memory	
Format:	EORIM Mx , I	
Instruction Code:	01100 iiii xxx xxxx	
Carry Flag:	Not affected	
Operation:	AC , $Mx \leftarrow Mx \wedge I$	

Explanation

Operand 2 of the instruction EORIM should be immediate data. This instruction is to do logic EOR op eration with the value of d ata memory and immediate data I, and to save the result in both the accumulator and the data memory.

[Example] \$20H=0 011B , I=0101B

1

LDI	20H,0011B	;\$20H=0011B,AC=0011B
EORIM	20H,0101B	;\$20H=0110B,AC=0110B
	:	



Programming Tip

After reading all of the logic inst ructions, you may wond er why there haven't the 'N OT' instruction (inverse)? If there haven't such an instruction in SH66 10 series, what can I do? Don't worry, Programming Tip is going to tell you how to use other instructions to perform the NOT function. The operation of NOT is to change $0 \rightarrow 1$ or $1 \rightarrow 0$ on each bit. Here, the EOR in struction can h elp us to g et the in versed value by doi ng EO R op eration wit h immediate data (0FH) and the variable that wants to be done NOT. You'll understand clearly after reading the following example.

[Example] Do the operation NOT with the value of \$20H (1100B)

EORIM 20H,0FH

Execution result: \$20H=0011B,AC=0011B

1-6 Instructions for Flow Control

Instruction: JMP	Function: to jump to a designated address to execute program
Format:	JMP X
Instruction Code:	1110p xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X(Include p)$

Explanation

Instruction JMP jump s to a desi gnated address to execute prog ram. Howeve r, addressing capability of SH6610 series CPU is limited to 4K wo rds (0000H~0FFFH), so the jumping range of JMP can only reach 4K(0FFFH). Address beyond 4K shall be reached by switching ba nks. There's detailed explanation in later ch apters for how to switch ban ks. The JMP instruction is similar to the GOTO instruction in BASIC program.

[Example] PC=40H, Jump to 0E00H

JMP 0E00H

Execution result: PC=0E00H



[Example] J ump to LABLE

:

JMP	LOOP	;jump to loop (PC ϵ	- 0340h)
	:		
ORG	0340H		
LOOP	: NOP		
		NOP	

Instruction: BAZ	Function: if AC equals 0, then go to a designated address to execute program; otherwise continue to execute the next line
Format:	BAZ X
Instruction Code:	10010 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X$, if $AC=0$

Explanation

If the value of AC is 0, then af ter executing the BAZ instruction, PC will go to the designated address X to execute p rogram, the range of X bein g from \$0 00H to \$7FF H. It continues to execute the next line if the value of AC is 1

[Example] if (\$20H=\$20H-1)=00H then goto INC21H

	LDI	20H,0FH	;\$20H=0FH
DEC20 H:	SBIM	20H,01H	;AC,\$20H ← \$20H -1
	BAZ	INC21H	;ifAC=0 jump to INC21H
	JMP	DEC20H	;else jump to <u>DEC20H</u>
		:	
INC21H	ADIM	21H,01H	;\$21H+1
		:	

Instruction: BA0	Function: if bit 0 of AC is 1, then go to a designated address to
	execute program; otherwise continue to execute the next
	line
Format:	BAO X
Instruction Code:	10100 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X$, if $AC(0)=1$



If bit 0 of AC is 1, then after executing BA0 instructio n, PC will go to the designated address X to execute program, the range of X being from \$000H to \$7FFH or from \$0800H to \$0FFFH. It continues to execute the next line if bit 0 of AC is 0.

[Example] if \$20H(bit 0)=1 then goto INC21H

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	LDI	20H,0FH	;\$20H=0FH,AC=0FH
DEC20H:	BIM	20H,01H	;\$20H,AC ← \$20H -1
	BA0	INC21H	;if AC(bit0)=1,jump to INC21H
	JMP	DEC20H	;else jump to DEC20H
		:	
INC21H:	ADIM	21H,01H	;\$21H,AC ← \$21H+1

Instruction: BA1	Function: if bit 1 of AC is 1, then go to designated address to execute program; otherwise continue to execute the next
	line
Format:	BA1 X
Instruction Code:	10101 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X$, if $AC(1)=1$

Explanation

If bit 1 of AC is 1, then after executing BA1 instructio n, PC will go to the designat ed address X to execute program, the range of X being from \$000H to \$7FFH or from \$0800H to \$0FFFH. It continues to execute the next line if bit 1 of AC is 0.

[Example] if \$20H(bit 1)=1 then goto INC21H

		•		
	LDI	20H,0FH	;\$20H=0FH,AC=	⊧0FH
DEC20 H	:	SBIM	20H,01H	;\$20H,AC ← \$20H -1
	BA1	INC21H	;ifAC(bit1)=1,jum	np to INC21H
		JMP	DEC20H	;else jump to DEC20H
		:		
INC21 H:	ADIM	21H,01H	;\$21H,AC ← \$21I	H+1
		:		



Instruction: BA2	Function: if bit 2 of AC is 1, then go to designated address to
	execute program; otherwise continue to execute the next line
Format:	BA2 X
Instruction Code:	10110 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X$, if $AC(2)=1$

Explanation

If bit 2 of AC is 1, then after executing BA2 instructio n, PC will go to the designat ed address X to execute program, the range of X being from \$000H to \$7FFH or from \$0800H or \$0FFFH. It continues to execute the next line if bit 1 of AC is 0.

[Example] if \$20H(bit 2)=1 then goto INC21H

:

	LDI	20H,0FH	;\$20H=0FH,AC=0FH
DEC20H:	SBIM	20H,01H	;\$20H,AC ← \$20H -1
	BA2	INC21H	;if AC(bit2)=1,jump to INC21H
	JMP	DEC20H	;else jump to DEC20H
		:	
INC21H:	ADIM	21H,01H	;\$21H,AC ← \$21H+1
		:	

Instruction: BA3	Function: if bit 3 of AC is 1, then go to designated address to execute program; otherwise continue to execute the next line
Format:	BA3 X
Instruction Code:	10111 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	$PC \leftarrow X$, if $AC(3)=1$

Explanation

If bit 3 of AC is 1, then after executing BA3 instruct ion, PC will go to the designated address X to execute program, the range of X being from \$000H to \$7FFH or from \$0800H to \$0FFFH. It continues to execute the next line if bit 1 of AC is 0.



[Example] if \$20H(bit 3)=1 then goto INC21H

		:	
	LDI	20H,0FH	;\$20H=0FH,AC=0FH
DEC20H	SBIM	20H,01H	;\$20H,AC ← \$20H -1
	BA3	INC21H	;if AC(bit3)=1,jump to INC21H
	JMP	DEC20H	;else jump to DEC20H
		:	
INC21H:	ADIM	21H,01H	;\$21H,AC ← \$21H+1
			:

Instruction: BC	Function: if CY is 1, then go to designated address to execute
	program; otherwise continue to execute the next line
Format:	BC X
Instruction Code:	10011 xxxx xxx xxxx
Carry Flag:	CY
Operation:	$PC \leftarrow X$, if $CY=1$

Explanation

If CY is 1, then after executing instruction BC, PC will go to the designated address X to execute program, the ran ge of X being from \$000 H to \$7FFH. It continues to execute the next line if CY is n ot 1. The instruction BC is often u sed after ad dition or subtraction to decide whether there is a carry or borrow. You should especially note that for addition, CY is set to 1 when there is a carry, while for subtraction CY is set to 1 when there isn't any borrow. Therefore you should be careful when dealing with program flows.

[Example] if CY=1 then goto INC20H

		•	
	LDI	20H,0FH	;\$20H=0FH,AC=0FH,CY=0
INC20H :	SBIM	20H,01H	;\$20H,AC ← \$20H -1
	BC	INC20H	;if CY=1,jump to INC20H
		:	



Instruction: TJMP	Function:Unconditionally go to the address composed by
	(PC11~PC8), TBR and AC value to execute program
Format:	TJMP
Instruction Code:	11110 1111 111 1111
Carry Flag:	Not affected
Operation:	

The destination address for Instruction TJMP is composed by PC's bit 8 ~bit 11, TBR and AC value (please refer to instruction RTNW for program example).

Example: PC=300H, TBR=01H, AC=02H, then the rule for composing a destination address is as follows:

address=PC(bit8~bit 11) TBR AC If: PC =<u>3</u>00H

TBR=01H

AC=02H

Then the destination address is: 3 1 2 H

Instruction: CALL	Function: to call a subprogram
Format:	CALL X
Instruction Code:	11000 xxxx xxx xxxx
Carry Flag:	Not affected
Operation:	ST \leftarrow CY; PC , PC \leftarrow X (not include p)

Explanation

Instruction CALL is used to call a subprogram. First it saves the values of CY and PC+1 to stack for returning to the calling program, then goes to the designated address X (\$0000H ~ \$07FFH or \$0800 H ~ \$0FFF H) to execute p rogram. Instructions RTNW or RTNI can be used to return to the calling program. When using CALL to call a subprogram, you sho uld especially note how many layers of st ack have al ready been used, because SH6610 series only provide 4-layer stacks. If more than 4 layers are used, serious error will be o ccurred when returning to the calling program!



Instruction: RTNW	Function: to return to the calling program, $H \rightarrow TBR$, $L \rightarrow AC$
Format:	RTNW H,L
Instruction Code:	11010 000h hhh 1111
Carry Flag:	Not affected
Operation:	PC \leftarrow ST , TBR \leftarrow hhhh , AC \leftarrow 1111

RTNW is an instruction to get data from stack to PC for returning to the calling program, and at the same time put the value of H into TBR and the value of L into AC. This instruction is often used to get stationary data.

[Example] To get data from ROM address 302H

TBR EQU		OEH	
TEMP	EQU	20H	
		:	
		:	
001A	LDI	TBR , 00H	;put index value (high nibble) 0 into TBR.
001B	LDI	TEMP , 02	;put index value (low nibble) 2 into AC
001C	CALL	300H	;call subprogram.
001D		:	
		:	
		:	
	ORG	300H	
0300	TJMP		; get destination address \$0302H according to
(PC11~PC8	B),TBR,AC		
0301	RTNW	00H,01H	
0302	RTNW	00H,02H	;return to main program, H→TBR,L→AC
0303	RTNW	04H,05H	
0304	RTNW	09H,08H	
0305		:	

Instruction: RTNI	Function: to return from interrupt or subprogram
Format:	RTNI
Instruction Code:	11010 1000 000 0000
Carry Flag:	CY
Operation:	CY;PC ←ST



Instruction RTNI is mainly used for returning from interrupt or sub program. It fills CY and PC with values of st ack (CY and returning add ress) when returning. What's the difference between RTNI and RTNW? We can find that when returning by RTNW, only the returning address in the st ack is fetched into PC, but CY value is not fetched. And RTNW fetches another two values (H \rightarrow TBR, L \rightarrow AC), which RTNI does not do. Therefore you can choose from the two instructions according to your needs.

[Example] To exchange two numbers

000E	1 TBR	EQU	0EH	
0020	2 REGX	EQU	20H	
0021	3 REGY	EQU	21H	
0022	4 TEMP	EQU	22H	
	5 ;***************			
0005 780E	12 RESET :	LDI	TEMP,00H	;set TEMP=00h
0006 7920	13	LDI	REGX,02H	;set RegX=02h
0007 7A21	14	LDI	REGY,04H	;set RegY=04h
0008 C00A	15	CALL	SWAPXY	; ;call subprogram
	17 ;*********************			
000A 3820	18 SWAPXY	LDA	REGX,00H	;AC=02H
000B 3C22	19	STA	TEMP,00H	;TEMP=02H
000C 3821	20	LDA	REGY,00H	;AC=04H
000D 3C20	21	STA	REGX,00H	;REGX=04H
000E 3822	22	LDA	TEMP,00H	;AC=02H
000F 3C21	23	STA	REGY,00H	;REG2=02H
0010 D400	24	RTNI	;return to	main program



Instruction: HALT	Function: CPU to be halt from working
Format:	HALT
Instruction Code:	11011 0000 000 0000
Carry Flag:	Not affected
Operation:	No

After executing inst ruction HALT, CPU will be ha It while it s surrounding circuit (counter, oscillation circuit) continues working. The instruct ion HALT is usually use d to stop CP U temporarily in order to save power. In HALT mode, when any of the system interrupts occurs, CPU will be released from HALT mode and continue to work.

[Example] HALT program, to enable PORT B interrupt to wake up the program

IEX	EQU	00H	;interrupt enable register
IRQ	EQU	01H	;interrupt require flag
PORT	BEQU	09H	;i/o port b
		:	
	LDI	PORTB,0FH	;set port b = " high '
	LDI	IEX,0001B	;enable port interrupt
	LDI	IRQ,00H	;clear interrupt require flag
	HALT		;system cpu halt
	NOP		
		:	

Instruction: STOP	Function: to stop the whole chip (including oscillation circuit)
Format:	STOP
Instruction Code:	11011 1000 000 0000
Carry Flag:	Not affected
Operation:	No

Explanation

Executing instruction STOP will stop the whole chip from working, including oscillation circuit. Only PORT interrupt and external interrupt can relea se CPU from STOP mode, so you must en able an interrupt before entering STOP mode, otherwise the system can't be waked up from STOP mode.



[Example] STOP program, to enable PORT B interrupt to wake up the program

IEX	EQU	00H	;interrupt enable register
IRQ	EQU	01H	;interrupt require flag
PORTB	EQU	09H	;i/o port b
		:	
	LDI	PORTB,0FH	;set port b = " high '
	LDI	IEX,0001B	;enable port interrupt
	LDI	IRQ,00H	;clear interrupt require flag
	STOP		;all system is " stop "
	NOP		
		:	

Instruction: NOP	Function: to do nothing
Format:	NOP
Instruction Code:	1111 1111 111 1111
Carry Flag:	Not affected
Operation:	No

Explanation

Instruction NOP means doing n othing in its instruction cycle and it is of ten used for time delay. Because it does n othing when executing, you don't worry if it will af fect any current status.