Features

- Single 2.5V 3.6V or 2.7V 3.6V Supply
- Serial Peripheral Interface (SPI) Compatible
- 20 MHz Max Clock Frequency
- Page Program Operation
 - Single Cycle Reprogram (Erase and Program)
- 4096 Pages (528 Bytes/Page) Main Memory
- Supports Page and Block Erase Operations
- Two 528-byte SRAM Data Buffers Allows Receiving of Data while Reprogramming of Nonvolatile Memory
- Continuous Read Capability through Entire Array
- Ideal for Code Shadowing Applications
- Low Power Dissipation
 - 4 mA Active Read Current Typical
 - 2 µA CMOS Standby Current Typical
- Hardware Data Protection Feature
- 100% Compatible to AT45DB161
- 5.0V-tolerant Inputs: SI, SCK, CS, RESET and WP Pins
- Commercial and Industrial Temperature Ranges
- Green (Pb/Halide-free) Packaging Options

Description

The AT45DB161B is a 2.5-volt or 2.7-volt only, serial interface Flash memory ideally suited for a wide variety of digital voice-, image-, program code- and data-storage

Pin Configurations

Function
Chip Select
Serial Clock
Serial Input
Serial Output
Hardware Page Write Protect Pin
Chip Reset
Ready/Busy

DataFlash Card⁽¹⁾

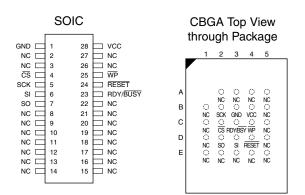
Top View through Package

1. See AT45DCB002 Datasheet.

Note:

TSOP Top View - Type 1

RDY/BUSY	1	28 🗖 NC
RESET	20	27 🗖 NC
WP	3	26 🗖 NC
NC 🗌	4	25 🗖 NC
NC 🗌	5	24 🗖 NC
VCC	6	23 🗖 NC
GND 🗔	7	22 🗖 NC
NC 🗆	8	21 🗖 NC
NC 🗌	9	20 🗖 NC
NC 🗌	10	19 🗖 NC
CS 🔤	11	18 🗖 NC
SCK 🗔	12	17 🗖 NC
SI	13	16 🗖 NC
SO 🗔	14	15 🗖 NC



CASON - Top View through Package

)
١C
С
P



16-megabit 2.5-volt Only or 2.7-volt Only DataFlash[®]

AT45DB161B

Rev. 2224I-DFLSH-10/04



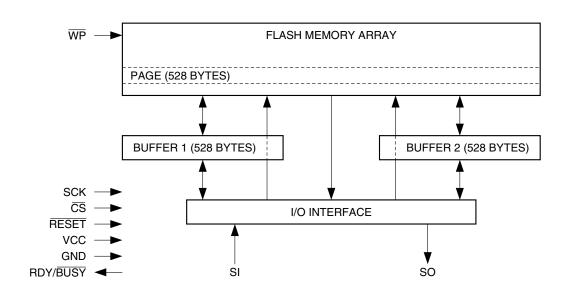


applications. Its 17,301,504 bits of memory are organized as 4096 pages of 528 bytes each. In addition to the main memory, the AT45DB161B also contains two SRAM data buffers of 528 bytes each. The buffers allow receiving of data while a page in the main memory is being reprogrammed, as well as writing a continuous data stream. EEPROM emulation (bit or byte alterability) is easily handled with a self-contained three step Read-Modify-Write operation.Unlike conventional Flash memories that are accessed randomly with multiple address lines and a parallel interface, the DataFlash uses a SPI serial interface to sequentially access its data. DataFlash supports SPI mode 0 and mode 3. The simple serial interface facilitates hardware layout, increases system reliability, minimizes switching noise, and reduces package size and active pin count. The device is optimized for use in many commercial and industrial applications where high density, low pin count, low voltage, and low power are essential. The device operates at clock frequencies up to 20 MHz with a typical active read current consumption of 4 mA.

To allow for simple in-system reprogrammability, the AT45DB161B does not require high input voltages for programming. The device operates from a single power supply, 2.5V to 3.6V or 2.7V to 3.6V, for both the program and read operations. The AT45DB161B is enabled through the chip select pin $\overline{(CS)}$ and accessed via a three-wire interface consisting of the Serial Input (SI), Serial Output (SO), and the Serial Clock (SCK).

All programming cycles are self-timed, and no separate erase cycle is required before programming.

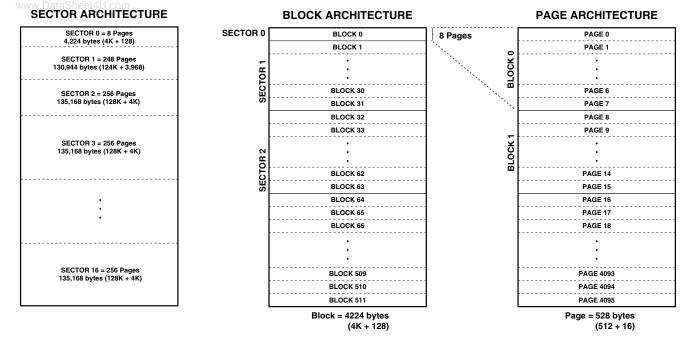
When the device is shipped from Atmel, the most significant page of the memory array may not be erased. In other words, the contents of the last page may not be filled with FFH.



Block Diagram

Memory Array

To provide optimal flexibility, the memory array of the AT45DB161B is divided into three levels of granularity comprising of sectors, blocks, and pages. The Memory Architecture Diagram illustrates the breakdown of each level and details the number of pages per sector and block. All program operations to the DataFlash occur on a page-by-page basis; however, the optional erase operations can be performed at the block or page level.



Memory Architecture Diagram

Device Operation

The device operation is controlled by instructions from the host processor. The list of instructions and their associated opcodes are contained in Tables 1 through 4. A valid instruction starts with the falling edge of \overline{CS} followed by the appropriate 8-bit opcode and the desired buffer or main memory address location. While the \overline{CS} pin is low, toggling the SCK pin controls the loading of the opcode and the desired buffer or main memory address location through the SI (serial input) pin. All instructions, addresses and data are transferred with the most significant bit (MSB) first.

Buffer addressing is referenced in the datasheet using the terminology BFA9 - BFA0 to denote the ten address bits required to designate a byte address within a buffer. Main memory addressing is referenced using the terminology PA11 - PA0 and BA9 - BA0 where PA11 - PA0 denotes the 12 address bits required to designate a page address and BA9 - BA0 denotes the ten address bits required to designate a byte address within the page.

Read Commands

By specifying the appropriate opcode, data can be read from the main memory or from either one of the two data buffers. The DataFlash supports two categories of read modes in relation to the SCK signal. The differences between the modes are in respect to the inactive state of the SCK signal as well as which clock cycle data will begin to be output. The two categories, which are comprised of four modes total, are defined as Inactive Clock Polarity Low or Inactive Clock Polarity High and SPI Mode 0 or SPI Mode 3. A separate opcode (refer to Table 1 on page 10 for a complete list) is used to select which category will be used for reading. Please refer to the "Detailed Bit-level Read Timing" diagrams in this datasheet for details on the clock cycle sequences for each mode.

CONTINUOUS ARRAY READ: By supplying an initial starting address for the main memory array, the Continuous Array Read command can be utilized to sequentially read a continuous stream of data from the device by simply providing a clock signal; no additional addressing information or control signals need to be provided. The DataFlash incorporates an internal address counter that will automatically increment on every clock





cycle, allowing one continuous read operation without the need of additional address sequences. To perform a continuous read, an opcode of 68H or E8H must be clocked into the device followed by 24 address bits and 32 don't care bits. The first two bits of the 24-bit address sequence are reserved for upward and downward compatibility to larger and smaller density devices (see Notes under "Command Sequence for Read/Write Operations" diagram). The next 12 address bits (PA11 - PA0) specify which page of the main memory array to read, and the last ten bits (BA9 - BA0) of the 24-bit address sequence specify the starting byte address within the page. The 32 don't care bits that follow the 24 address bits are needed to initialize the read operation. Following the 32 don't care bits, additional clock pulses on the SCK pin will result in serial data being output on the SO (serial output) pin.

The CS pin must remain low during the loading of the opcode, the address bits, the don't care bits, and the reading of data. When the end of a page in main memory is reached during a Continuous Array Read, the device will continue reading at the beginning of the next page with no delays incurred during the page boundary crossover (the crossover from the end of one page to the beginning of the next page). When the last bit in the main memory array has been read, the device will continue reading back at the beginning of the first page of memory. As with crossing over page boundaries, no delays will be incurred when wrapping around from the end of the array to the beginning of the array.

A low-to-high transition on the \overline{CS} pin will terminate the read operation and tri-state the SO pin. The maximum SCK frequency allowable for the Continuous Array Read is defined by the f_{CAR} specification. The Continuous Array Read bypasses both data buffers and leaves the contents of the buffers unchanged.

MAIN MEMORY PAGE READ: A Main Memory Page Read allows the user to read data directly from any one of the 4096 pages in the main memory, bypassing both of the data buffers and leaving the contents of the buffers unchanged. To start a page read, an opcode of 52H or D2H must be clocked into the device followed by 24 address bits and 32 don't care bits. The first two bits of the 24-bit address sequence are reserved bits, the next 12 address bits (PA11 - PA0) specify the page address, and the next ten address bits (BA9 - BA0) specify the starting byte address within the page. The 32 don't care bits which follow the 24 address bits are sent to initialize the read operation. Following the 32 don't care bits, additional pulses on SCK result in serial data being output on the SO (serial output) pin. The \overline{CS} pin must remain low during the loading of the opcode, the address bits, the don't care bits, and the reading of data. When the end of a page in main memory is reached during a Main Memory Page Read, the device will continue reading at the beginning of the same page. A low-to-high transition on the \overline{CS} pin will terminate the read operation and tri-state the SO pin.

BUFFER READ: Data can be read from either one of the two buffers, using different opcodes to specify which buffer to read from. An opcode of 54H or D4H is used to read data from buffer 1, and an opcode of 56H or D6H is used to read data from buffer 2. To perform a Buffer Read, the eight bits of the opcode must be followed by 14 don't care bits, ten address bits, and eight don't care bits. Since the buffer size is 528 bytes, ten address bits (BFA9 - BFA0) are required to specify the first byte of data to be read from the buffer. The \overline{CS} pin must remain low during the loading of the opcode, the address bits, the don't care bits, and the reading of data. When the end of a buffer is reached, the device will continue reading back at the beginning of the buffer. A low-to-high transition on the \overline{CS} pin will terminate the read operation and tri-state the SO pin.

STATUS REGISTER READ: The status register can be used to determine the device's Ready/Busy status, the result of a Main Memory Page to Buffer Compare operation, or the device density. To read the status register, an opcode of 57H or D7H must be

www.DataShe

loaded into the device. After the last bit of the opcode is shifted in, the eight bits of the status register, starting with the MSB (bit 7), will be shifted out on the SO pin during the next eight clock cycles. The five most significant bits of the status register will contain device information, while the remaining three least-significant bits are reserved for future use and will have undefined values. After bit 0 of the status register has been shifted out, the sequence will repeat itself (as long as \overline{CS} remains low and SCK is being tog-gled) starting again with bit 7. The data in the status register is constantly updated, so each repeating sequence will output new data.

Status Register Format

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
RDY/BUSY	COMP	1	0	1	1	Х	Х

Ready/Busy status is indicated using bit 7 of the status register. If bit 7 is a 1, then the device is not busy and is ready to accept the next command. If bit 7 is a 0, then the device is in a busy state. The user can continuously poll bit 7 of the status register by stopping SCK at a low level once bit 7 has been output. The status of bit 7 will continue to be output on the SO pin, and once the device is no longer busy, the state of SO will change from 0 to 1. There are eight operations which can cause the device to be in a busy state: Main Memory Page to Buffer Transfer, Main Memory Page to Buffer Compare, Buffer to Main Memory Page Program with Built-in Erase, Buffer to Main Memory Page Program without Built-in Erase, Page Erase, Block Erase, Main Memory Page Program, and Auto Page Rewrite.

The result of the most recent Main Memory Page to Buffer Compare operation is indicated using bit 6 of the status register. If bit 6 is a 0, then the data in the main memory page matches the data in the buffer. If bit 6 is a 1, then at least one bit of the data in the main memory page does not match the data in the buffer.

The device density is indicated using bits 5, 4, 3 and 2 of the status register. For the AT45DB161B, the four bits are 1, 0, 1 and 1. The decimal value of these four binary bits does not equate to the device density; the four bits represent a combinational code relating to differing densities of Serial DataFlash devices, allowing a total of sixteen different density configurations.

Program and Erase
CommandsBUFFER WRITE: Data can be shifted in from the SI pin into either buffer 1 or buffer 2.
To load data into either buffer, an 8-bit opcode, 84H for buffer 1 or 87H for buffer 2, must
be followed by 14 don't care bits and ten address bits (BFA9 - BFA0). The ten address
bits specify the first byte in the buffer to be written. The data is entered following the
address bits. If the end of the data buffer is reached, the device will wrap around back to
the beginning of the buffer. Data will continue to be loaded into the buffer until a low-to-
high transition is detected on the CS pin.

BUFFER TO MAIN MEMORY PAGE PROGRAM WITH BUILT-IN ERASE: Data written into either buffer 1 or buffer 2 can be programmed into the main memory. To start the operation, an 8-bit opcode, 83H for buffer 1 or 86H for buffer 2, must be followed by the two reserved bits, 12 address bits (PA11 - PA0) that specify the page in the main memory to be written, and ten additional don't care bits. When a low-to-high transition occurs on the \overline{CS} pin, the part will first erase the selected page in main memory to all 1s and then program the data stored in the buffer into the specified page in the main memory. Both the erase and the programming of the page are internally self-timed and should take place in a maximum time of t_{EP}. During this time, the status register will indicate that the part is busy.





BUFFER TO MAIN MEMORY PAGE PROGRAM WITHOUT BUILT-IN ERASE: A

www.DataShe

previously erased page within main memory can be programmed with the contents of either buffer 1 or buffer 2. To start the operation, an 8-bit opcode, 88H for buffer 1 or 89H for buffer 2, must be followed by the two reserved bits, 12 address bits (PA11 - PA0) that specify the page in the main memory to be written, and ten additional don't care bits. When a low-to-high transition occurs on the \overline{CS} pin, the part will program the data stored in the buffer into the specified page in the main memory. It is necessary that the page in main memory that is being programmed has been previously erased. The programming of the page is internally self-timed and should take place in a maximum time of t_P. During this time, the status register will indicate that the part is busy.

Successive page programming operations without doing a page erase are not recommended. In other words, changing bytes within a page from a "1" to a "0" during multiple page programming operations without erasing that page is not recommended.

PAGE ERASE: The optional Page Erase command can be used to individually erase any page in the main memory array allowing the Buffer to Main Memory Page Program without Built-in Erase command to be utilized at a later time. To perform a Page Erase, an opcode of 81H must be loaded into the device, followed by two reserved bits, 12 address bits (PA11 - PA0), and ten don't care bits. The 12 address bits are used to specify which page of the memory array is to be erased. When a low-to-high transition occurs on the \overline{CS} pin, the part will erase the selected page to 1s. The erase operation is internally self-timed and should take place in a maximum time of t_{PE} . During this time, the status register will indicate that the part is busy.

BLOCK ERASE: A block of eight pages can be erased at one time allowing the Buffer to Main Memory Page Program without Built-in Erase command to be utilized to reduce programming times when writing large amounts of data to the device. To perform a Block Erase, an opcode of 50H must be loaded into the device, followed by two reserved bits, nine address bits (PA11 - PA3), and 13 don't care bits. The nine address bits are used to specify which block of eight pages is to be erased. When a low-to-high transition occurs on the \overline{CS} pin, the part will erase the selected block of eight pages to 1s. The erase operation is internally self-timed and should take place in a maximum time of t_{BE}. During this time, the status register will indicate that the part is busy.

PA11	PA10	PA9	PA8	PA7	PA6	PA5	PA4	PA3	PA2	PA1	PA0	Block
0	0	0	0	0	0	0	0	0	Х	Х	Х	0
0	0	0	0	0	0	0	0	1	х	х	х	1
0	0	0	0	0	0	0	1	0	х	х	х	2
0	0	0	0	0	0	0	1	1	х	х	х	3
•	•	•	•	•	•	•	•	•	•	•	•	•
•	•	•	•	•	•	•	•	•	•	•	•	•
•	•	•	•	•	•	•	•	•	•	•	•	•
1	1	1	1	1	1	1	0	0	х	х	Х	508
1	1	1	1	1	1	1	0	1	х	х	х	509
1	1	1	1	1	1	1	1	0	х	х	х	510
1	1	1	1	1	1	1	1	1	х	х	х	511

Block Erase Addressing

MAIN MEMORY PAGE PROGRAM THROUGH BUFFER: This operation is a combination of the Buffer Write and Buffer to Main Memory Page Program with Built-in Erase operations. Data is first shifted into buffer 1 or buffer 2 from the SI pin and then programmed into a specified page in the main memory. To initiate the operation, an 8-bit opcode, 82H for buffer 1 or 85H for buffer 2, must be followed by the two reserved bits and 22 address bits. The 12 most significant address bits (PA11 - PA0) select the page in the main memory where data is to be written, and the next ten address bits (BFA9 - BFA0) select the first byte in the buffer to be written. After all address bits are shifted in, the part will take data from the SI pin and store it in one of the data buffers. If the end of the buffer is reached, the device will wrap around back to the beginning of the buffer. When there is a low-to-high transition on the \overline{CS} pin, the part will first erase the selected page in main memory to all 1s and then program the data stored in the buffer into the specified page in the main memory. Both the erase and the programming of the page are internally self-timed and should take place in a maximum of time t_{EP}. During this time, the status register will indicate that the part is busy.

Additional Commands MAIN MEMORY PAGE TO BUFFER TRANSFER: A page of data can be transferred from the main memory to either buffer 1 or buffer 2. To start the operation, an 8-bit opcode, 53H for buffer 1 and 55H for buffer 2, must be followed by the two reserved bits, 12 address bits (PA11 - PA0) which specify the page in main memory that is to be transferred, and ten don't care bits. The CS pin must be low while toggling the SCK pin to load the opcode, the address bits, and the don't care bits from the SI pin. The transfer of the page of data from the main memory to the buffer will begin when the CS pin transitions from a low to a high state. During the transfer of a page of data (t_{XFR}), the status register can be read to determine whether the transfer has been completed or not.

MAIN MEMORY PAGE TO BUFFER COMPARE: A page of data in main memory can be compared to the data in buffer 1 or buffer 2. To initiate the operation, an 8-bit opcode, 60H for buffer 1 and 61H for buffer 2, must be followed by 24 address bits consisting of the two reserved bits, 12 address bits (PA11 - PA0) which specify the page in the main memory that is to be compared to the buffer, and ten don't care bits. The \overline{CS} pin must be low while toggling the SCK pin to load the opcode, the address bits, and the don't care bits from the SI pin. On the low-to-high transition of the \overline{CS} pin, the 528 bytes in the selected main memory page will be compared with the 528 bytes in buffer 1 or buffer 2. During this time (t_{XFR}), the status register will indicate that the part is busy. On completion of the compare operation, bit 6 of the status register is updated with the result of the compare.

AUTO PAGE REWRITE: This mode is only needed if multiple bytes within a page or multiple pages of data are modified in a random fashion. This mode is a combination of two operations: Main Memory Page to Buffer Transfer and Buffer to Main Memory Page Program with Built-in Erase. A page of data is first transferred from the main memory to buffer 1 or buffer 2, and then the same data (from buffer 1 or buffer 2) is programmed back into its original page of main memory. To start the rewrite operation, an 8-bit opcode, 58H for buffer 1 or 59H for buffer 2, must be followed by the two reserved bits, 12 address bits (PA11 - PA0) that specify the page in main memory to be rewritten, and ten additional don't care bits. When a low-to-high transition occurs on the \overline{CS} pin, the part will first transfer data from the page in main memory. The operation is internally self-timed and should take place in a maximum time of t_{EP}. During this time, the status register will indicate that the part is busy.





Operation Mode Summary

If a sector is programmed or reprogrammed sequentially page-by-page, then the programming algorithm shown in Figure 1 on page 26 is recommended. Otherwise, if multiple bytes in a page or several pages are programmed randomly in a sector, then the programming algorithm shown in Figure 2 on page 27 is recommended. Each page within a sector must be updated/rewritten at least once within every 10,000 cumulative page erase/program operations in that sector.

The modes described can be separated into two groups – modes which make use of the Flash memory array (Group A) and modes which do not make use of the Flash memory array (Group B).

Group A modes consist of:

- 1. Main Memory Page Read
- 2. Main Memory Page to Buffer 1 (or 2) Transfer
- 3. Main Memory Page to Buffer 1 (or 2) Compare
- 4. Buffer 1 (or 2) to Main Memory Page Program with Built-in Erase
- 5. Buffer 1 (or 2) to Main Memory Page Program without Built-in Erase
- 6. Page Erase
- 7. Block Erase
- 8. Main Memory Page Program through Buffer
- 9. Auto Page Rewrite

Group B modes consist of:

- 1. Buffer 1 (or 2) Read
- 2. Buffer 1 (or 2) Write
- 3. Status Register Read

If a Group A mode is in progress (not fully completed) then another mode in Group A should not be started. However, during this time in which a Group A mode is in progress, modes in Group B can be started.

This gives the Serial DataFlash the ability to virtually accommodate a continuous data stream. While data is being programmed into main memory from buffer 1, data can be loaded into buffer 2 (or vice versa). See application note AN-4 ("Using Atmel's Serial DataFlash") for more details.

Pin Descriptions SERIAL INPUT (SI): The SI pin is an input-only pin and is used to shift data into the device. The SI pin is used for all data input including opcodes and address sequences.

SERIAL OUTPUT (SO): The SO pin is an output-only pin and is used to shift data out from the device.

SERIAL CLOCK (SCK): The SCK pin is an input-only pin and is used to control the flow of data to and from the DataFlash. Data is always clocked into the device on the rising edge of SCK and clocked out of the device on the falling edge of SCK.

CHIP SELECT (\overline{CS}): The DataFlash is selected when the \overline{CS} pin is low. When the device is not selected, data will not be accepted on the SI pin, and the SO pin will remain in a high-impedance state. A high-to-low transition on the \overline{CS} pin is required to start an operation, and a low-to-high transition on the \overline{CS} pin is required to end an operation.

WRITE PROTECT: If the $\overline{\text{WP}}$ pin is held low, the first 256 pages of the main memory cannot be reprogrammed. The only way to reprogram the first 256 pages is to first drive the protect pin high and then use the program commands previously mentioned. If this pin and feature are not utilized it is recommended that the $\overline{\text{WP}}$ pin be driven high externally.

RESET: A low state on the reset pin (RESET) will terminate the operation in progress and reset the internal state machine to an idle state. The device will remain in the reset condition as long as a low level is present on the RESET pin. Normal operation can resume once the RESET pin is brought back to a high level.

The device incorporates an internal power-on reset circuit, so there are no restrictions on the RESET pin during power-on sequences. If this pin and feature are not utilized it is recommended that the RESET pin be driven high externally.

READY/BUSY: This open drain output pin will be driven low when the device is busy in an internally self-timed operation. This pin, which is normally in a high state (through a 1 k Ω external pull-up resistor), will be pulled low during programming operations, compare operations, and during page-to-buffer transfers.

The busy status indicates that the Flash memory array and one of the buffers cannot be accessed; read and write operations to the other buffer can still be performed.

Power-on/Reset State

When power is first applied to the device, or when recovering from a reset condition, the device will default to SPI Mode 3. In addition, the SO pin will be in a high-impedance state, and a high-to-low transition on the \overline{CS} pin will be required to start a valid instruction. The SPI mode will be automatically selected on every falling edge of \overline{CS} by sampling the inactive clock state. After power is applied and V_{CC} is at the minimum datasheet value, the system should wait 20 ms before an operational mode is started.





Table 1. Read Commands

Command	SCK Mode	Opcode			
Continuous Array Dood	Inactive Clock Polarity Low or High	68H			
Continuous Array Read	SPI Mode 0 or 3	E8H			
Main Memory Page Read Buffer 1 Read Buffer 2 Read	Inactive Clock Polarity Low or High	52H			
	SPI Mode 0 or 3	D2H			
	Inactive Clock Polarity Low or High	54H			
	SPI Mode 0 or 3	D4H			
	Inactive Clock Polarity Low or High	56H			
	SPI Mode 0 or 3	D6H			
Ototuo Dogistar Dood	Inactive Clock Polarity Low or High	57H			
Status Register Read	SPI Mode 0 or 3	D7H			

Table 2. Program and Erase Commands

Command	SCK Mode	Opcode
Buffer 1 Write	Any	84H
Buffer 2 Write	Any	87H
Buffer 1 to Main Memory Page Program with Built-in Erase	Any	83H
Buffer 2 to Main Memory Page Program with Built-in Erase	Any	86H
Buffer 1 to Main Memory Page Program without Built-in Erase	Any	88H
Buffer 2 to Main Memory Page Program without Built-in Erase	Any	89H
Page Erase	Any	81H
Block Erase	Any	50H
Main Memory Page Program through Buffer 1	Any	82H
Main Memory Page Program through Buffer 2	Any	85H

Table 3. Additional Commands

Command	SCK Mode	Opcode
Main Memory Page to Buffer 1 Transfer	Any	53H
Main Memory Page to Buffer 2 Transfer	Any	55H
Main Memory Page to Buffer 1 Compare	Any	60H
Main Memory Page to Buffer 2 Compare	Any	61H
Auto Page Rewrite through Buffer 1	Any	58H
Auto Page Rewrite through Buffer 2	Any	59H

Note: In Tables 2 and 3, an SCK mode designation of "Any" denotes any one of the four modes of operation (Inactive Clock Polarity Low, Inactive Clock Polarity High, SPI Mode 0, or SPI Mode 3).

	Table 4.	Detailed Bit-level Addressing Sequence
--	----------	--

www.	Address Byte						Address Byte								Address Byte											
Opcode	Opcode	Reserved	Reserved	PA11	PA10	PA9	PA8	PA7	PA6	PA5	PA4	PA3	PA2	PA1	PA0	BA9	BA8	BA7	BA6	BA5	BA4	BA3	BA2	BA1	BA0	Additional Don't Care Bytes Required
50H	0 1 0 1 0 0 0 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	х	х	х	x	х	х	х	х	х	х	х	N/A
52H	0 1 0 1 0 0 1 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	В	в	в	в	в	4 Bytes
53H	01010011	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	x	N/A
54H	0 1 0 1 0 1 0 0	x	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	1 Byte
55H	01010101	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	х	х	х	х	х	х	х	х	N/A
56H	0 1 0 1 0 1 1 0	x	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	1 Byte
57H	01010111				Ν	J/A							1	J/A							١	J/A				N/A
58H	0 1 0 1 1 0 0 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
59H	01011001	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
60H	0 1 1 0 0 0 0 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
61H	0 1 1 0 0 0 0 1	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
68H	0 1 1 0 1 0 0 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	в	в	в	в	в	4 Bytes
81H	10000001	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
82H	10000010	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	в	в	в	в	в	N/A
83H	10000011	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
84H	10000100	x	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	N/A
85H	10000101	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	в	в	в	в	в	N/A
86H	10000110	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
87H	10000111	x	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	N/A
88H	10001000	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
89H	10001001	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	х	х	x	х	х	х	х	х	х	х	N/A
D2H	1 1 0 1 0 0 1 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	в	в	в	в	в	4 Bytes
D4H	1 1 0 1 0 1 0 0	х	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	1 Byte
D6H	1 1 0 1 0 1 1 0	x	х	х	х	х	х	х	х	x	х	х	х	х	х	в	в	в	в	в	в	в	в	в	в	1 Byte
D7H	1 1 0 1 0 1 1 1				Ν	J/A							1	۱/A							١	I/A				N/A
E8H	1 1 1 0 1 0 0 0	r	r	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	Ρ	в	в	в	в	в	в	в	в	в	в	4 Bytes

Note: r = Reserved Bit

P = Page Address Bit B = Byte/Buffer Address Bit

x = Don't Care





Absolute Maximum Ratings*

Temperature under Bias55°C to +125°C
Storage Temperature
All Input Voltages (including NC Pins) with Respect to Ground0.6V to +6.25V
All Output Voltages with Respect to Ground0.6V to V _{CC} + 0.6V

*NOTICE: Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

DC and AC Operating Range

		AT45DB161B (2.5V Version)	AT45DB161B
Operating Temperature (Case)	Com.	0°C to 70°C	0°C to 70°C
Operating Temperature (Case)	Ind.	_	-40°C to 85°C
V _{CC} Power Supply ⁽¹⁾		2.5V to 3.6V	2.7V to 3.6V

Note: 1. After power is applied and V_{CC} is at the minimum specified datasheet value, the system should wait 20 ms before an operational mode is started.

DC Characteristics

Symbol	Parameter	ameter Condition		Тур	Max	Units	
I _{SB}	Standby Current	\overline{CS} , \overline{RESET} , $\overline{WP} = V_{CC}$, all inputs at CMOS levels		2	10	μA	
I _{CC1} ⁽¹⁾	Active Current, Read $f = 20 \text{ MHz}; I_{OUT} = 0 \text{ mA};$ Operation $V_{CC} = 3.6V$			4	10	mA	
I _{CC2}	Active Current, Program/Erase Operation	V _{CC} = 3.6V		15	35	mA	
ILI	Input Load Current	V _{IN} = CMOS levels			1	μA	
I _{LO}	Output Leakage Current V _{I/O} = CMOS levels				1	μA	
V _{IL}	Input Low Voltage				0.6	V	
V _{IH}	Input High Voltage		2.0			V	
V _{OL}	Output Low Voltage	I _{OL} = 1.6 mA; V _{CC} = 2.7V			0.4	V	
V _{OH}	Output High Voltage	I _{OH} = -100 μA	V _{CC} - 0.2V			V	

Note: 1. I_{cc1} during a buffer read is 20mA maximum.

AC Characteristics

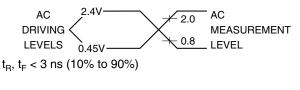
		AT45DB161B (2.5V Version)		AT45DB161B		
Symbol	Parameter	Min	Max	Min Max		Units
f _{scк}	SCK Frequency		15		20	MHz
f _{CAR}	SCK Frequency for Continuous Array Read		15		20	MHz
t _{WH}	SCK High Time	30		22		ns
t _{WL}	SCK Low Time	30		22		ns
t _{CS}	Minimum CS High Time	250		250		ns
t _{CSS}	CS Setup Time	250		250		ns
t _{CSH}	CS Hold Time	250		250		ns
t _{CSB}	CS High to RDY/BUSY Low		200		200	ns
t _{su}	Data In Setup Time	10		5		ns
t _H	Data In Hold Time	15		10		ns
t _{HO}	Output Hold Time	0		0		ns
t _{DIS}	Output Disable Time		20		18	ns
t _v	Output Valid		25		20	ns
t _{XFR}	Page to Buffer Transfer/Compare Time		300		250	μs
t _{EP}	Page Erase and Programming Time		20		20	ms
t _P	Page Programming Time		14		14	ms
t _{PE}	Page Erase Time		8		8	ms
t _{BE}	Block Erase Time		12		12	ms
t _{RST}	RESET Pulse Width	10		10		μs
t _{REC}	RESET Recovery Time		1		1	μs



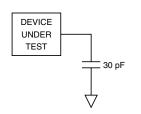


Input Test Waveforms and Measurement Levels





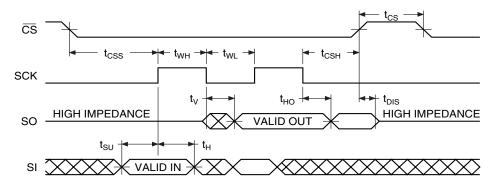
Output Test Load



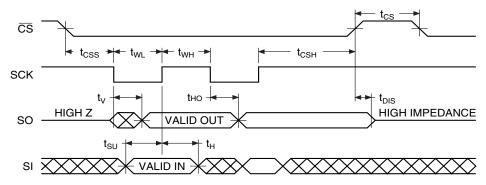
AC Waveforms Two different timing diagrams are shown below. Waveform 1 shows the SCK signal being low when \overline{CS} makes a high-to-low transition, and Waveform 2 shows the SCK signal being high when \overline{CS} makes a high-to-low transition. Both waveforms show valid timing diagrams. The setup and hold times for the SI signal are referenced to the low-to-high transition on the SCK signal.

Waveform 1 shows timing that is also compatible with SPI Mode 0, and Waveform 2 shows timing that is compatible with SPI Mode 3.

Waveform 1 – Inactive Clock Polarity Low and SPI Mode 0

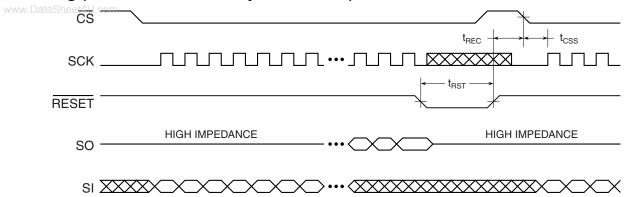


Waveform 2 – Inactive Clock Polarity High and SPI Mode 3



14 **AT45DB161B**

Reset Timing (Inactive Clock Polarity Low Shown)



Note: The CS signal should be in the high state before the RESET signal is deasserted.

s X

Command Sequence for Read/Write Operations (except Status Register Read)

CMD



8 bits

8 bits

8 bits

	T	
Reserved for	Page Address	Byte/Buffer Address
larger densities	(PA11-PA0)	(BA9-BA0/BFA9-BFA0)

Notes: 1. "r" designates bits reserved for larger densities.

- 2. It is recommended that "r" be a logical "0" for densities of 16M bits or smaller.
- 3. For densities larger than 16M bits, the "r" bits become the most significant Page Address bit for the appropriate density.

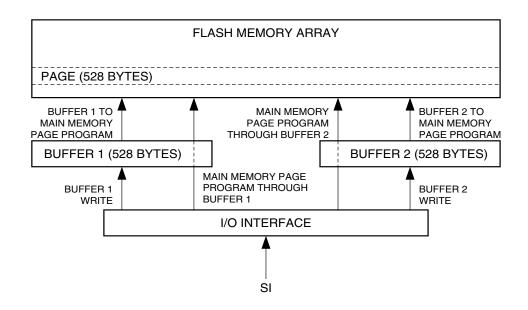




Write Operations

www.DataSheet4U.com

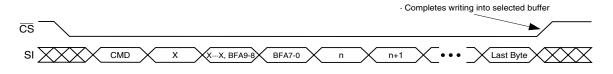
The following block diagram and waveforms illustrate the various write sequences available.



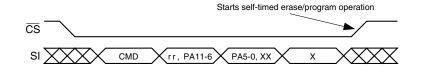
Main Memory Page Program through Buffers

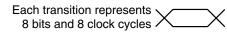


Buffer Write



Buffer to Main Memory Page Program (Data from Buffer Programmed into Flash Page)



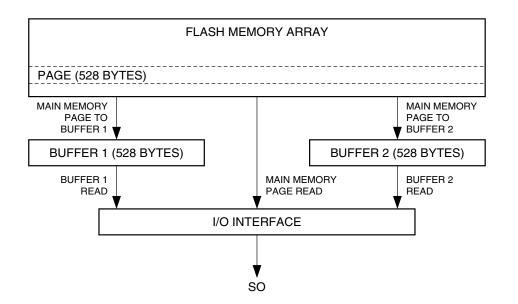


n = 1st byte read n+1 = 2nd byte read

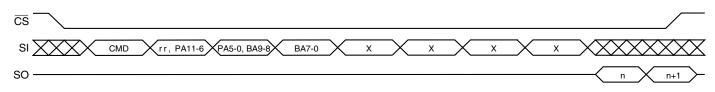
Read Operations

www.DataSheet4U.com

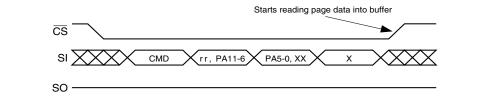
The following block diagram and waveforms illustrate the various read sequences available.



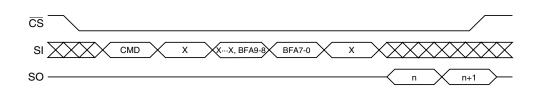
Main Memory Page Read



Main Memory Page to Buffer Transfer (Data from Flash Page Read into Buffer)



Buffer Read



n = 1st byte read n+1 = 2nd byte read

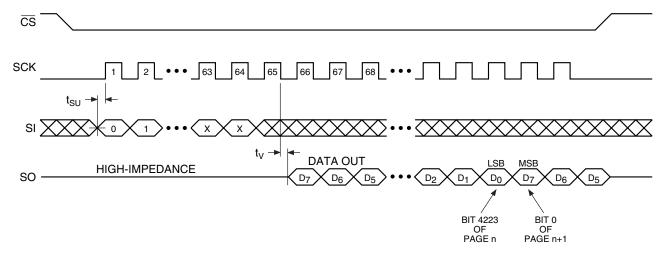




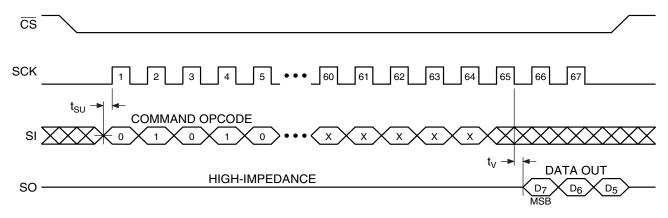
Detailed Bit-level Read Timing – Inactive Clock Polarity Low

/ww.DataSheet4U.com

Continuous Array Read (Opcode: 68H)



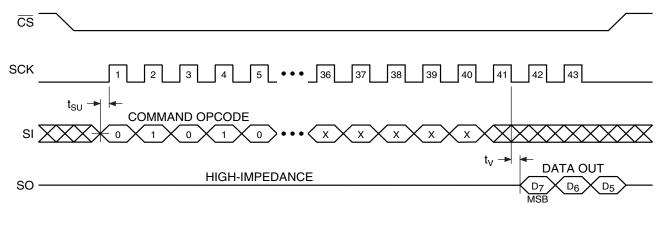
Main Memory Page Read (Opcode: 52H)



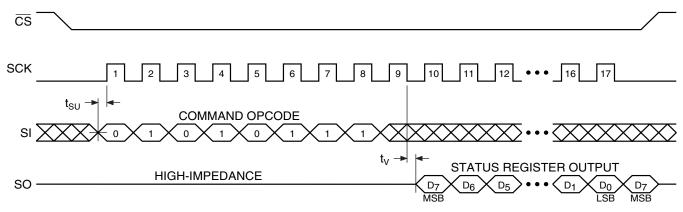
Detailed Bit-level Read Timing – Inactive Clock Polarity Low (Continued)

ww.DataSheet4U.com

Buffer Read (Opcode: 54H or 56H)



Status Register Read (Opcode: 57H)



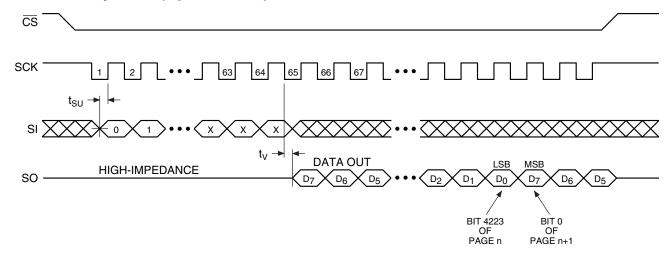




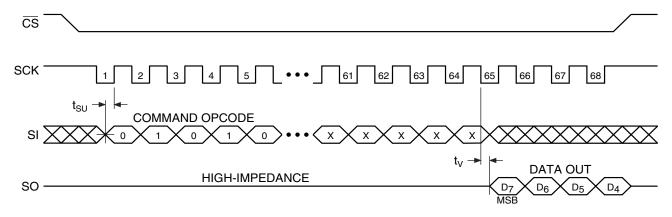
Detailed Bit-level Read Timing – Inactive Clock Polarity High

/ww.DataSheet4U.com

Continuous Array Read (Opcode: 68H)



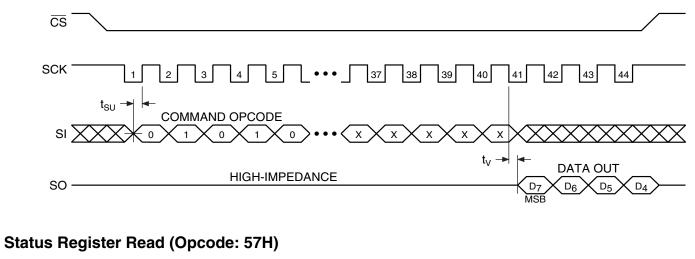
Main Memory Page Read (Opcode: 52H)

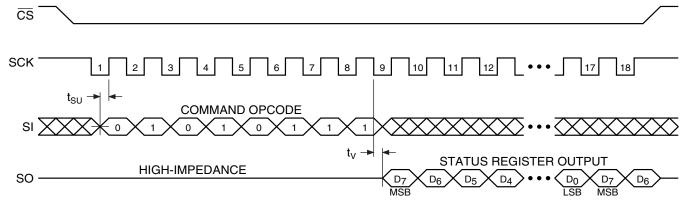


Detailed Bit-level Read Timing – Inactive Clock Polarity High (Continued)

ww.DataSheet4U.com

Buffer Read (Opcode: 54H or 56H)





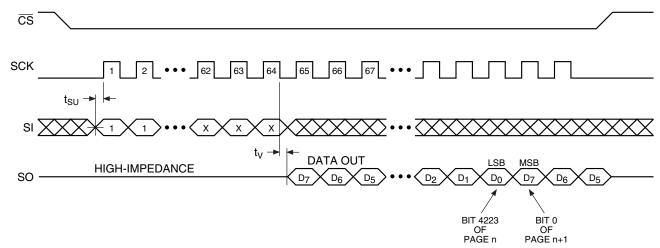




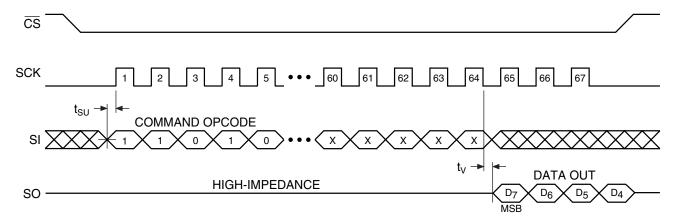
Detailed Bit-level Read Timing – SPI Mode 0

ww.DataSheet4U.com

Continuous Array Read (Opcode: E8H)



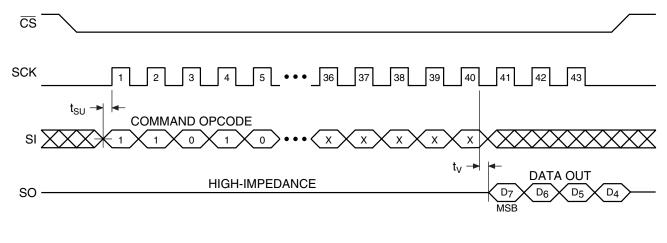
Main Memory Page Read (Opcode: D2H)



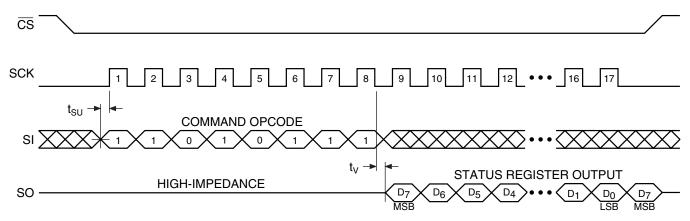
Detailed Bit-level Read Timing – SPI Mode 0 (Continued)

ww.DataSheet4U.com

Buffer Read (Opcode: D4H or D6H)



Status Register Read (Opcode: D7H)



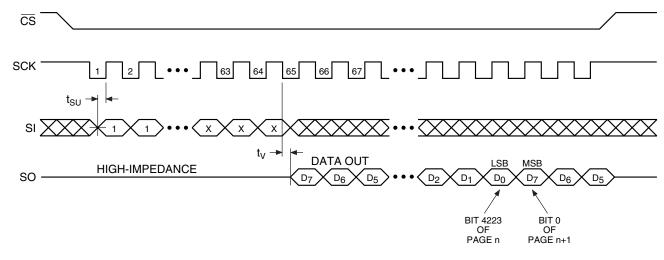




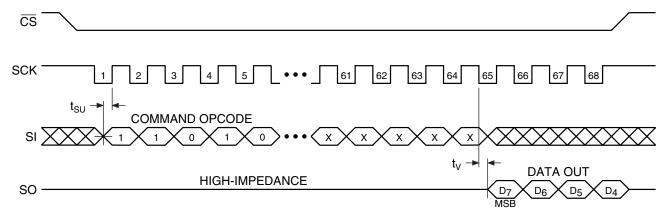
Detailed Bit-level Read Timing – SPI Mode 3

vww.DataSheet4U.com

Continuous Array Read (Opcode: E8H)



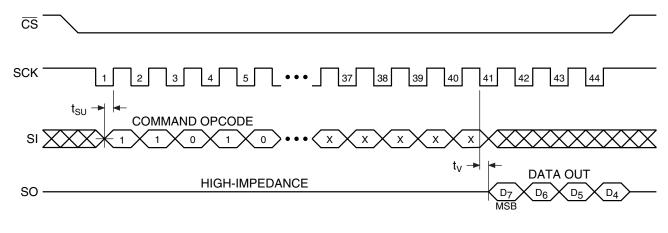
Main Memory Page Read (Opcode: D2H)



Detailed Bit-level Read Timing – SPI Mode 3 (Continued)

ww.DataSheet4U.com

Buffer Read (Opcode: D4H or D6H)



Status Register Read (Opcode: D7H)

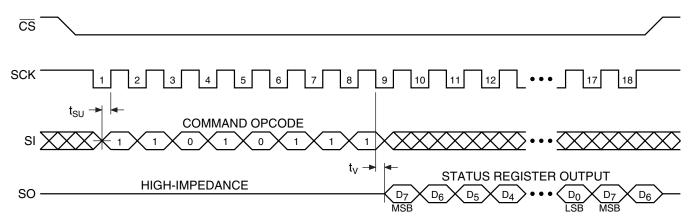
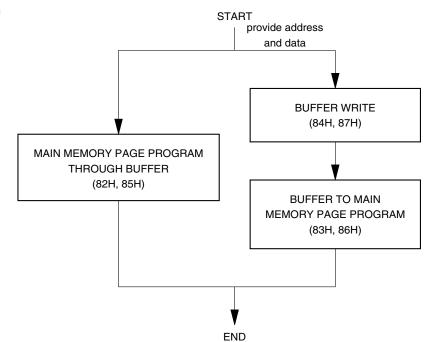






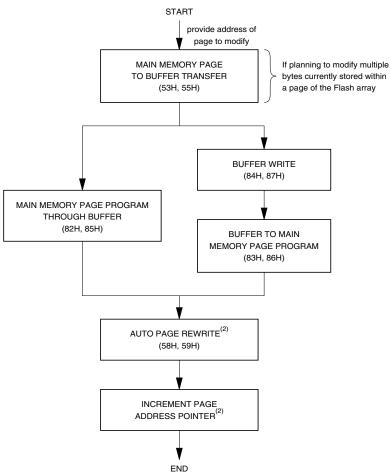
Figure 1. Algorithm for Sequentially Programming or Reprogramming the Entire Array



- Notes: 1. This type of algorithm is used for applications in which the entire array is programmed sequentially, filling the array page-bypage.
 - 2. A page can be written using either a Main Memory Page Program operation or a Buffer Write operation followed by a Buffer to Main Memory Page Program operation.
 - 3. The algorithm above shows the programming of a single page. The algorithm will be repeated sequentially for each page within the entire array.

Figure 2. Algorithm for Randomly Modifying Data

www.DataSheet4U.cor



- Notes: 1. To preserve data integrity, each page of a DataFlash sector must be updated/rewritten at least once within every 10,000 cumulative page erase/program operations.
 - 2. A Page Address Pointer must be maintained to indicate which page is to be rewritten. The Auto Page Rewrite command must use the address specified by the Page Address Pointer.
 - 3. Other algorithms can be used to rewrite portions of the Flash array. Low-power applications may choose to wait until 10,000 cumulative page erase/program operations have accumulated before rewriting all pages of the sector. See application note AN-4 ("Using Atmel's Serial DataFlash") for more details.

PA11	PA10	PA9	PA8	PA7	PA6	PA5	PA4	PA3	Sector
0	0	0	0	0	0	0	0	0	0
0	0	0	0	Х	Х	Х	Х	Х	1
0	0	0	1	Х	Х	Х	Х	х	2
0	0	1	0	Х	Х	Х	Х	Х	3
•	•	•	•	•	•	•	•	•	•
•	•	•	•	•	•	•	•	•	•
•	•	•	•	•	•	•	•	•	•
1	1	0	0	Х	Х	Х	Х	Х	13
1	1	0	1	Х	Х	Х	Х	х	14
1	1	1	0	Х	Х	Х	Х	Х	15
1	1	1	1	Х	Х	Х	Х	Х	16

Sector Addressing





Ordering Information

f _{scк}	I _{cc}	(mA)					
(MHz)	Active	Standby	Ordering Code	Package	Operation Range		
15	10	0.01	AT45DB161B-CC-2.5 AT45DB161B-CNC-2.5 AT45DB161B-RC-2.5 AT45DB161B-RC-2.5	24C1 8CN3 28R 28T	Commercial (0°C to 70°C) 2.5V to 3.6V		
20	10	0.01	AT45DB161B-CC AT45DB161B-CNC AT45DB161B-RC AT45DB161B-TC	24C1 8CN3 28R 28T	Commercial (0°C to 70°C)		
20	10	0.01	AT45DB161B-CI AT45DB161B-CNI AT45DB161B-RI AT45DB161B-TI	24C1 8CN3 28R 28T	Industrial (-40°C to 85°C)		

Green Packaging Options (Pb/Halide-free)

f _{sck}	f _{scк} I _{cc} (mA)				
(MHz)	Active	Standby	Ordering Code	Package	Operation Range
20	10	0.01	AT45DB161B-CNU AT45DB161B-RU AT45DB161B-TU	8CN3 28R 28T	Industrial (-40°C to 85°C)

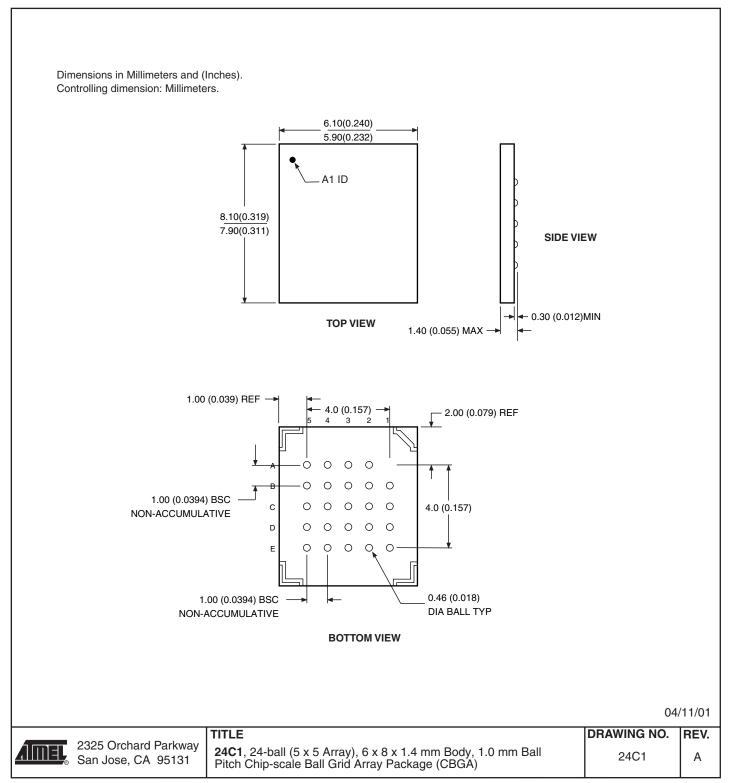
Note: Green packages cover lead-free requirements.

	Package Type						
24C1	24-ball (5 x 5 Array), Plastic Chip-scale Ball Grid Array (CBGA)						
8CN3	8-pad (6 mm x 8 mm) Chip Array Small Outline No Lead Package (CASON)						
28R	28-lead, 0.330" Wide, Plastic Gull Wing Small Outline Package (SOIC)						
28T	28-lead, Plastic Thin Small Outline Package (TSOP)						

Packaging Information

ww.DataSheet4U.com

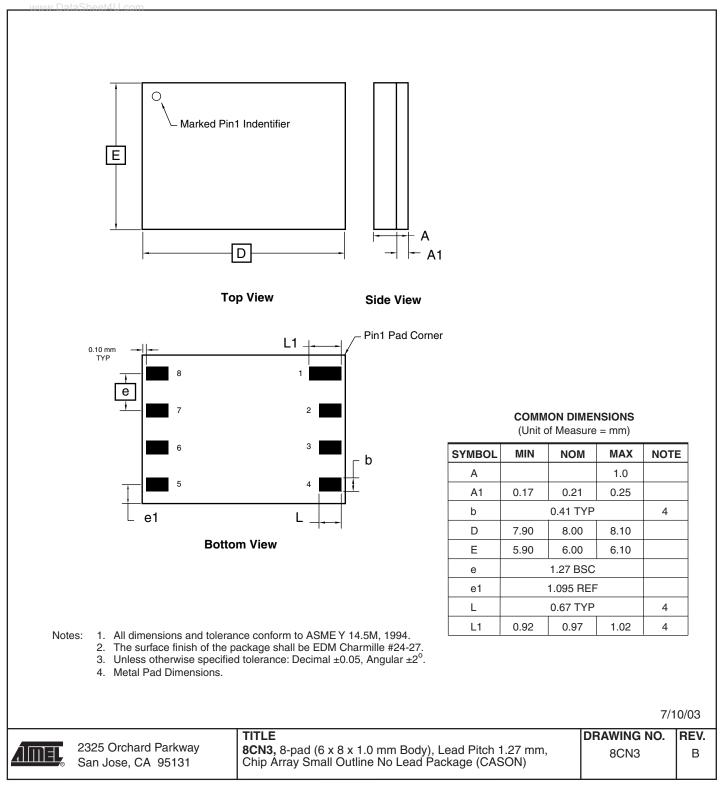
24C1 – CBGA



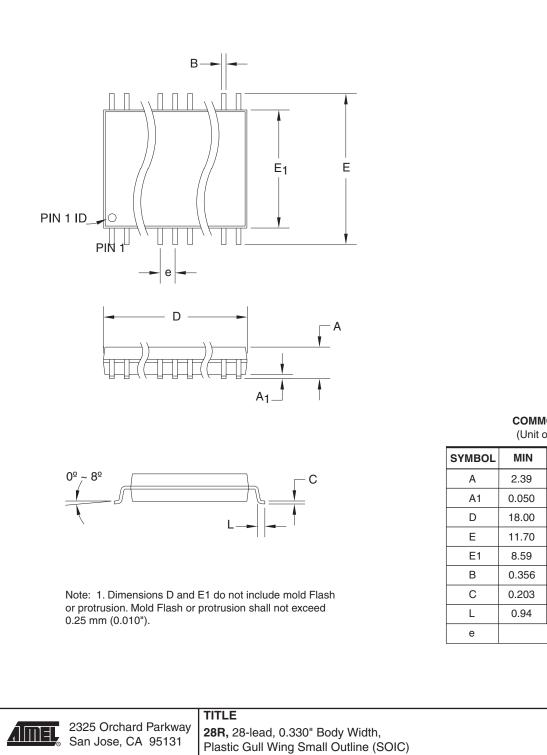




8CN3 – CASON



28R – SOIC





(Unit of Measure = mm)

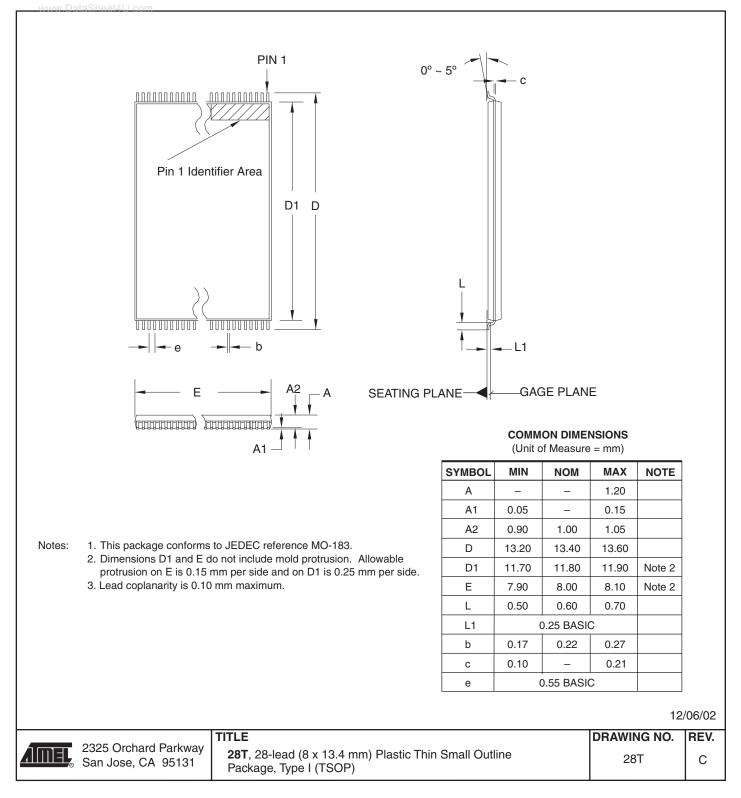
SYMBOL	MIN	NOM	MAX	NOTE
А	2.39	_	2.79	
A1	0.050	_	0.356	
D	18.00	-	18.50	Note 1
Е	11.70	-	12.50	
E1	8.59	-	8.79	Note 1
В	0.356	-	0.508	
С	0.203	-	0.305	
L	0.94	-	1.27	
е				

		5/18/2	2004
0005 Ouch and Darkman	TITLE	DRAWING NO.	REV.
2325 Orchard Parkway San Jose, CA 95131	28R, 28-lead, 0.330" Body Width, Plastic Gull Wing Small Outline (SOIC)	28R	с





28T – TSOP





Atmel Corporation

2325 Orchard Parkway San Jose, CA 95131, USA Tel: 1(408) 441-0311 Fax: 1(408) 487-2600

Regional Headquarters

Europe

Atmel Sarl Route des Arsenaux 41 Case Postale 80 CH-1705 Fribourg Switzerland Tel: (41) 26-426-5555 Fax: (41) 26-426-5500

Asia

Room 1219 Chinachem Golden Plaza 77 Mody Road Tsimshatsui East Kowloon Hong Kong Tel: (852) 2721-9778 Fax: (852) 2722-1369

Japan

9F, Tonetsu Shinkawa Bldg. 1-24-8 Shinkawa Chuo-ku, Tokyo 104-0033 Japan Tel: (81) 3-3523-3551 Fax: (81) 3-3523-7581

Atmel Operations

Memory

2325 Orchard Parkway San Jose, CA 95131, USA Tel: 1(408) 441-0311 Fax: 1(408) 436-4314

Microcontrollers

2325 Orchard Parkway San Jose, CA 95131, USA Tel: 1(408) 441-0311 Fax: 1(408) 436-4314

La Chantrerie BP 70602 44306 Nantes Cedex 3, France Tel: (33) 2-40-18-18-18 Fax: (33) 2-40-18-19-60

ASIC/ASSP/Smart Cards

Zone Industrielle 13106 Rousset Cedex, France Tel: (33) 4-42-53-60-00 Fax: (33) 4-42-53-60-01

1150 East Cheyenne Mtn. Blvd. Colorado Springs, CO 80906, USA Tel: 1(719) 576-3300 Fax: 1(719) 540-1759

Scottish Enterprise Technology Park Maxwell Building East Kilbride G75 0QR, Scotland Tel: (44) 1355-803-000 Fax: (44) 1355-242-743

RF/Automotive

Theresienstrasse 2 Postfach 3535 74025 Heilbronn, Germany Tel: (49) 71-31-67-0 Fax: (49) 71-31-67-2340

1150 East Cheyenne Mtn. Blvd. Colorado Springs, CO 80906, USA Tel: 1(719) 576-3300 Fax: 1(719) 540-1759

Biometrics/Imaging/Hi-Rel MPU/

High Speed Converters/RF Datacom Avenue de Rochepleine BP 123 38521 Saint-Egreve Cedex, France Tel: (33) 4-76-58-30-00 Fax: (33) 4-76-58-34-80

Literature Requests www.atmel.com/literature

Disclaimer: The information in this document is provided in connection with Atmel products. No license, express or implied, by estoppel or otherwise, to any intellectual property right is granted by this document or in connection with the sale of Atmel products. EXCEPT AS SET FORTH IN ATMEL'S TERMS AND CONDI-TIONS OF SALE LOCATED ON ATMEL'S WEB SITE, ATMEL ASSUMES NO LIABILITY WHATSOEVER AND DISCLAIMS ANY EXPRESS, IMPLIED OR STATUTORY WARRANTY RELATING TO ITS PRODUCTS INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTY OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE, OR NON-INFRINGEMENT. IN NO EVENT SHALL ATMEL BE LIABLE FOR ANY DIRECT, INDIRECT, CONSEQUENTIAL, PUNTIVE, SPECIAL OR INCIDEN-TAL DAMAGES (INCLUDING, WITHOUT LIMITATION, DAMAGES FOR LOSS OF PROFITS, BUSINESS INTERRUPTION, OR LOSS OF INFORMATION) ARISING OUT OF THE USE OR INABILITY TO USE THIS DOCUMENT, EVEN IF ATMEL HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES. Atmel makes no representations or warranties with respect to the accuracy or completeness of the contents of this document and reserves the right to make changes to specifications and product descriptions at any time without notice. Atmel does not make any commitment to update the information contained herein. Atmel's products are not intended, authorized, or warranted for use as components in applications intended to support or sustain life.

© Atmel Corporation 2004. All rights reserved. Atmel[®], logo and combinations thereof, and DataFlash[®] are registered trademarks, and Everywhere You AreSM is the trademark of Atmel Corporation or its subsidiaries. Other terms and product names may be trademarks of others.

